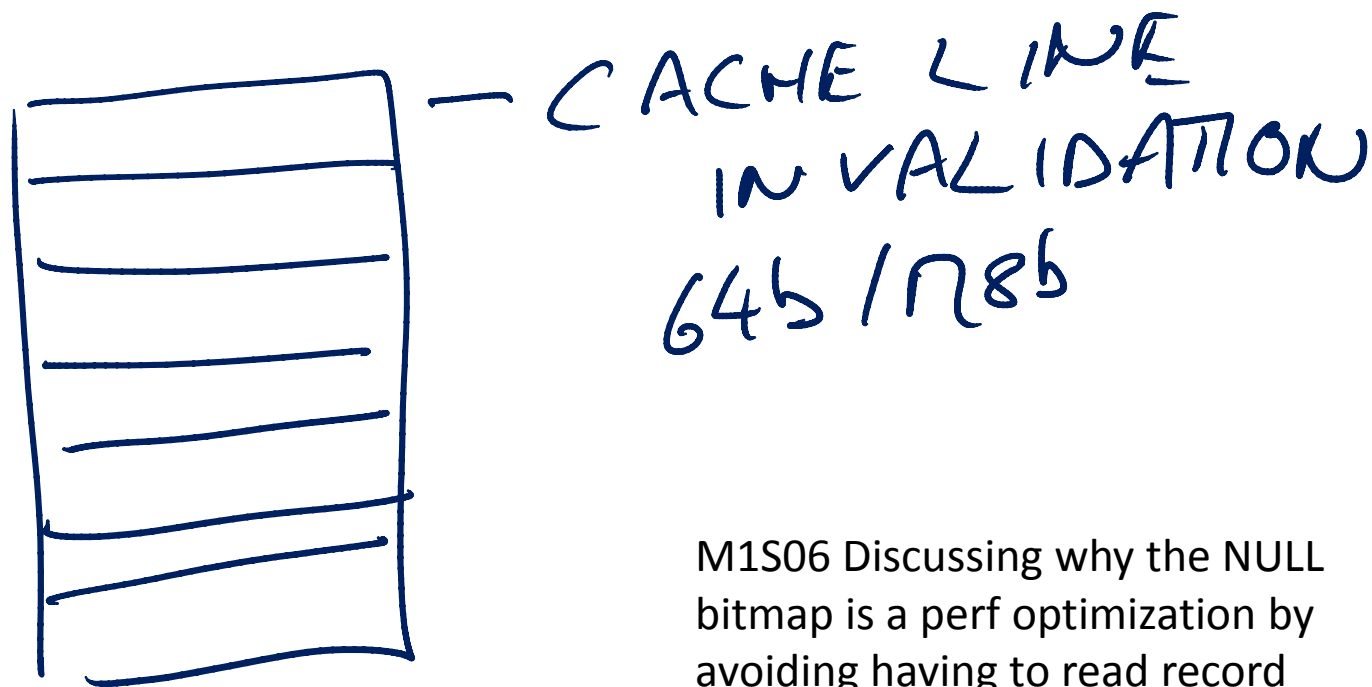
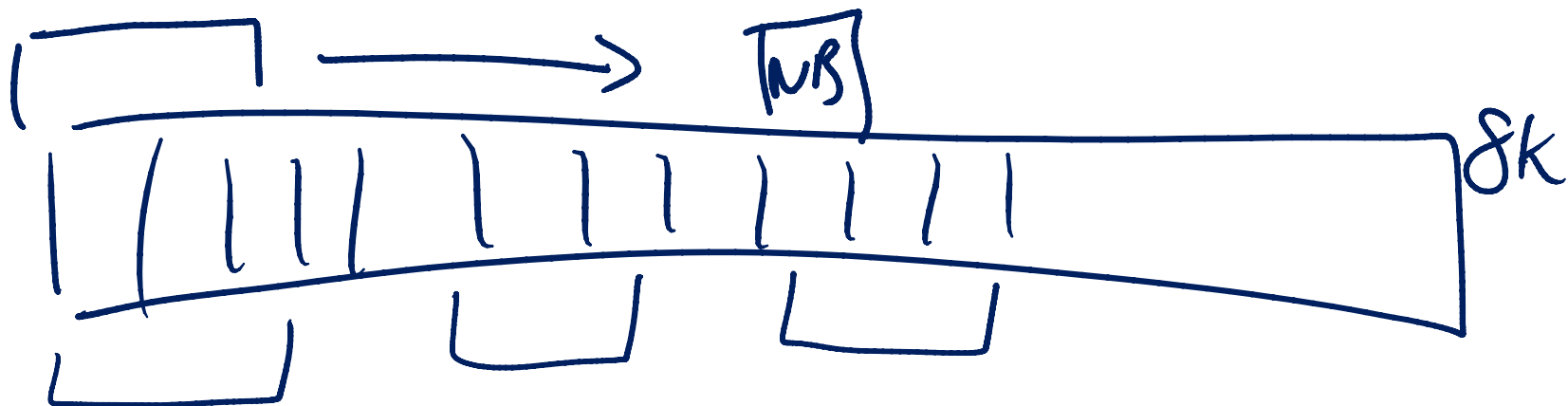
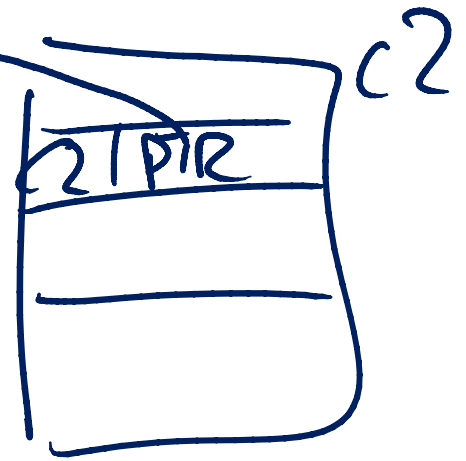
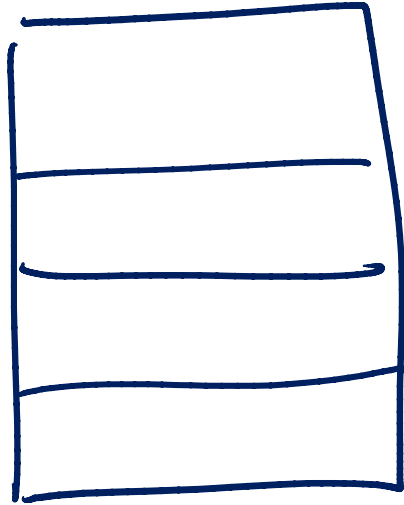




M1S06 Discussing why the NULL bitmap is a perf optimization by avoiding having to read record parts into the CPU's L1 cache

HEAP $c_1, \dots, c_{10}$

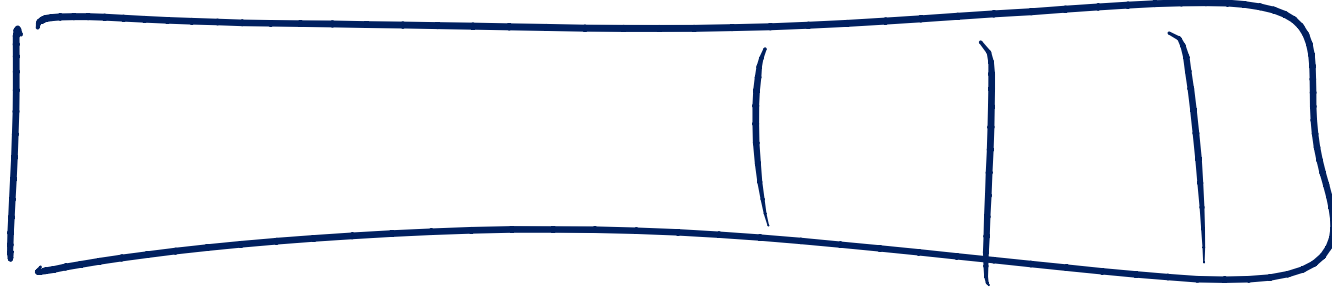


sel c_1, c_2, c_3
from table
where $c_2 = x$

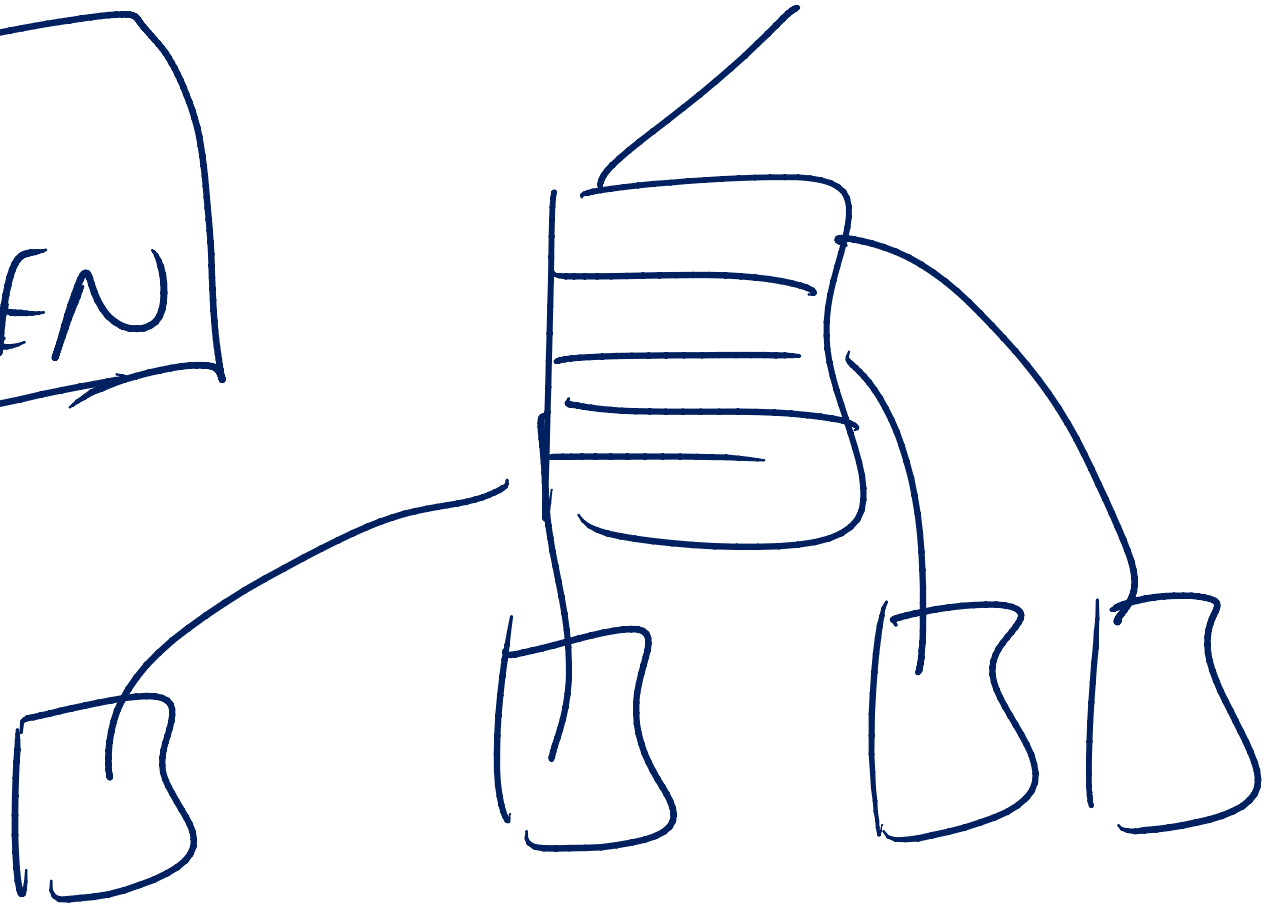
PTR	
HEAP	PHYSICAL RID (FILE:PAGE:SLOT)
c_2	LOGICAL RID (CLUSTER KEY)

M1S09 Explaining NC index back-links to the data records as part of describing why forwarding/forwarded records exist. Without them, moving a heap row would entail changing all NC index records with a back-link to that row.

M1S11 Off-row LOBs are stored in a loose tree structure. Google Marc Rasmussen for blog post with internals. I don't cover as there's nothing you can change/influence.



MARC
RASMUSSEN

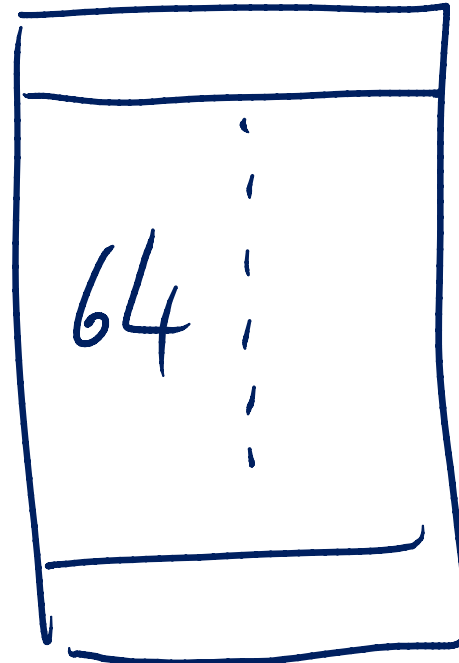
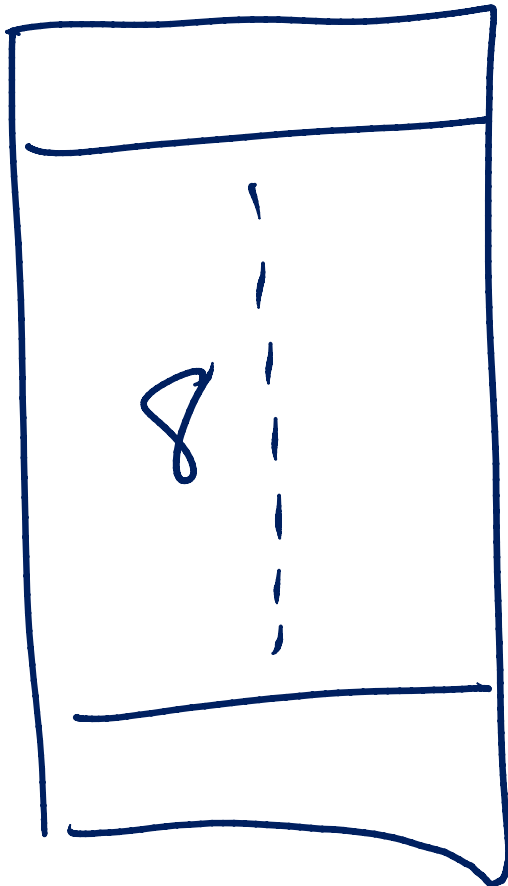


1K

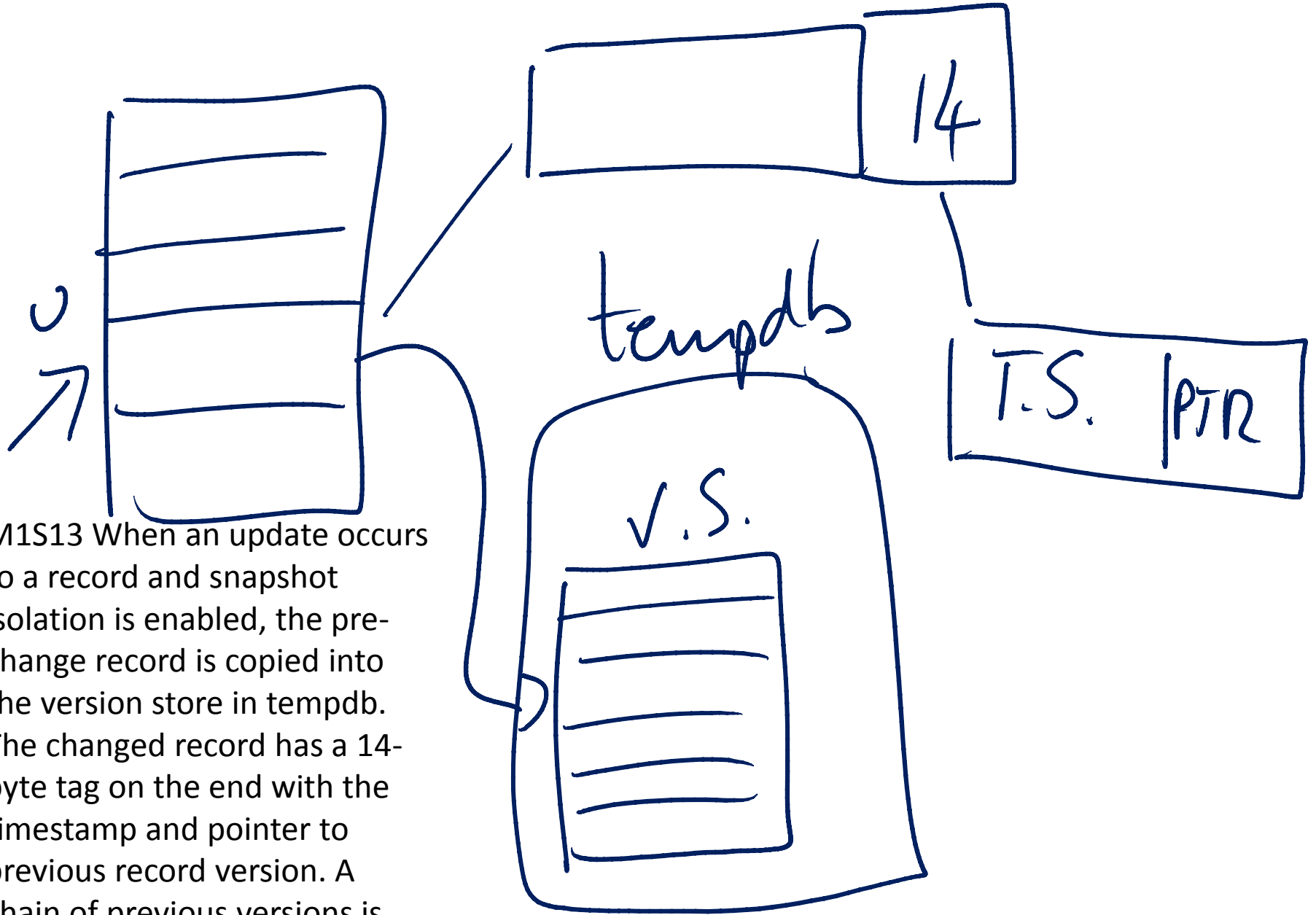
100	900	1%
-----	-----	----

100	24
-----	----

900

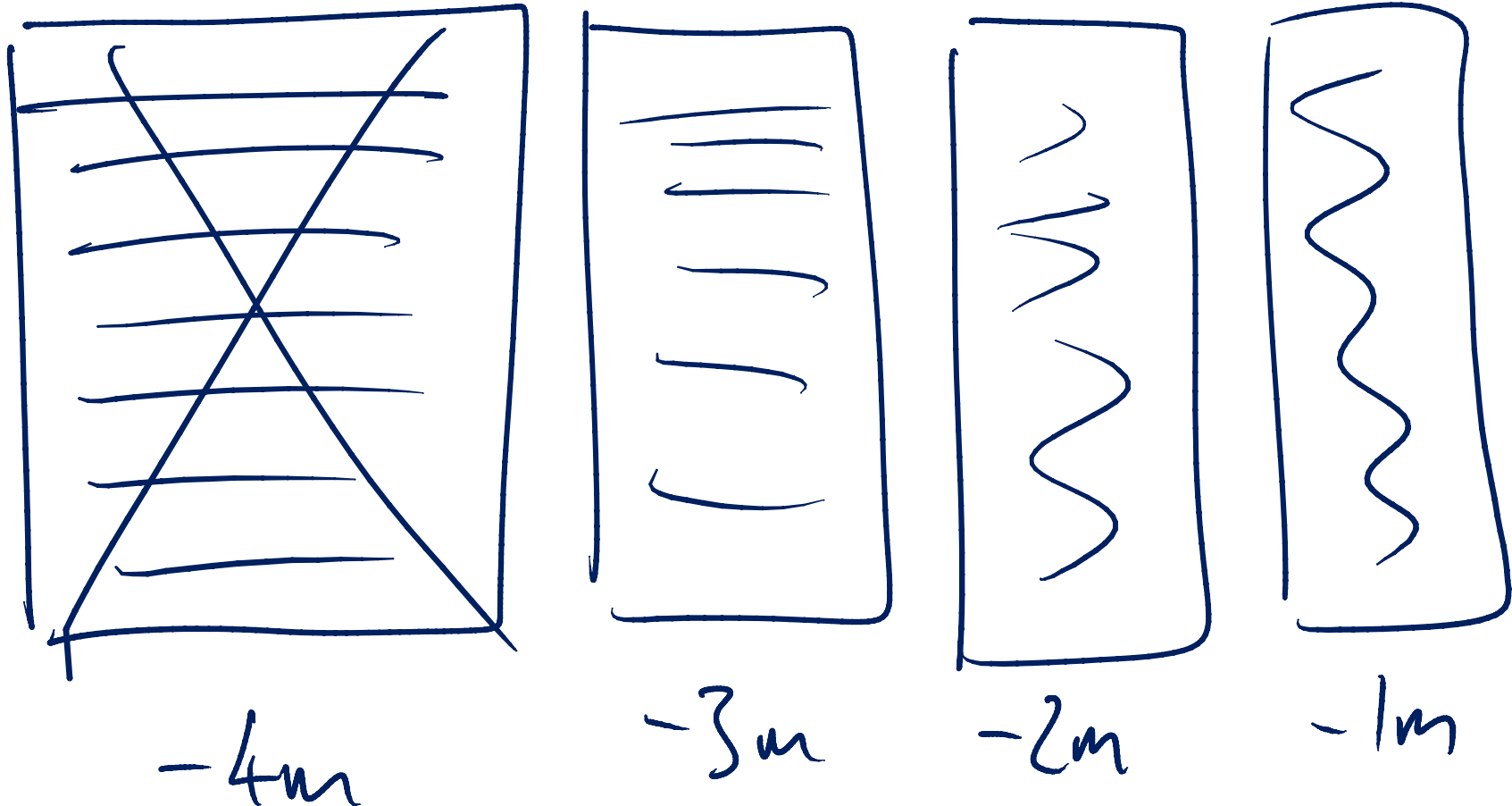


M1S11 Having a large LOB value in row that's not used much makes for large rows and low data density. Choose how to store LOB data wisely, either in-row, off-row, or in another table with a join.



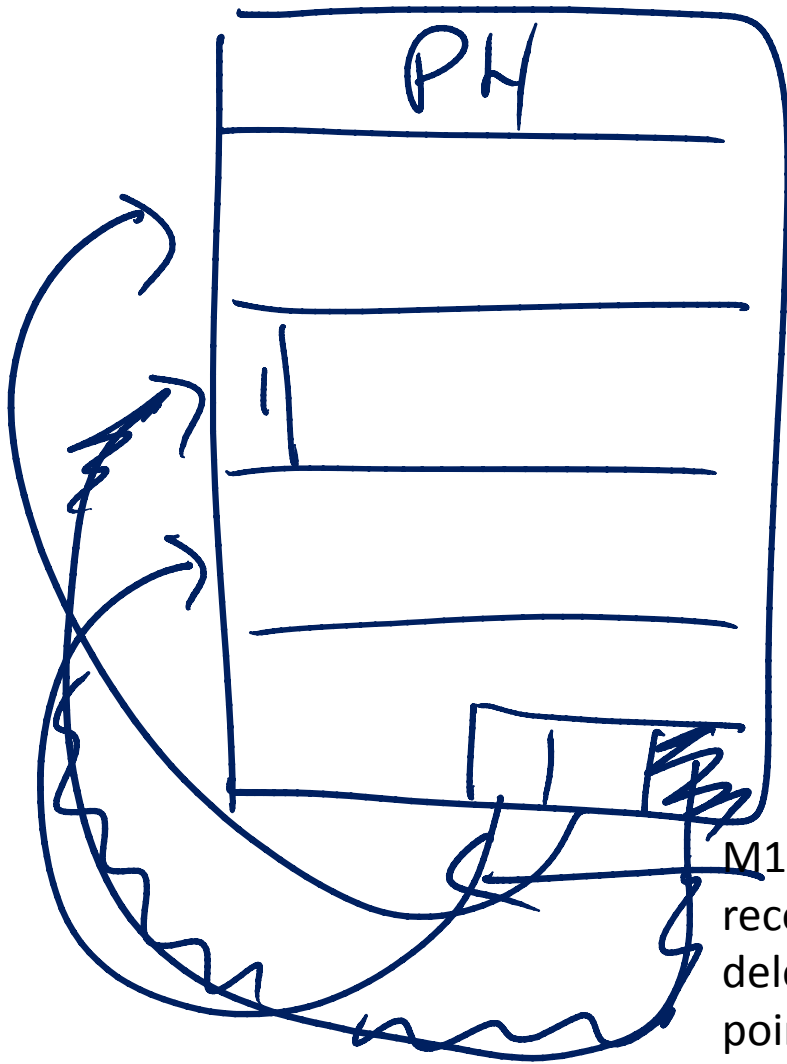
M1S13 When an update occurs to a record and snapshot isolation is enabled, the pre-change record is copied into the version store in tempdb. The changed record has a 14-byte tag on the end with the timestamp and pointer to previous record version. A chain of previous versions is possible.

APPEND-ONLY ALLOCATION UNIT



M1S13 and M2 S46 When an update occurs to a record and snapshot isolation is enabled, the pre-change record is copied into the version store in tempdb. The changed record has a 14-byte tag on the end with the timestamp and pointer to previous record version. A chain of previous versions is possible.

LOP_EXPUNCE
- GHOST



SLOT
ARRAY

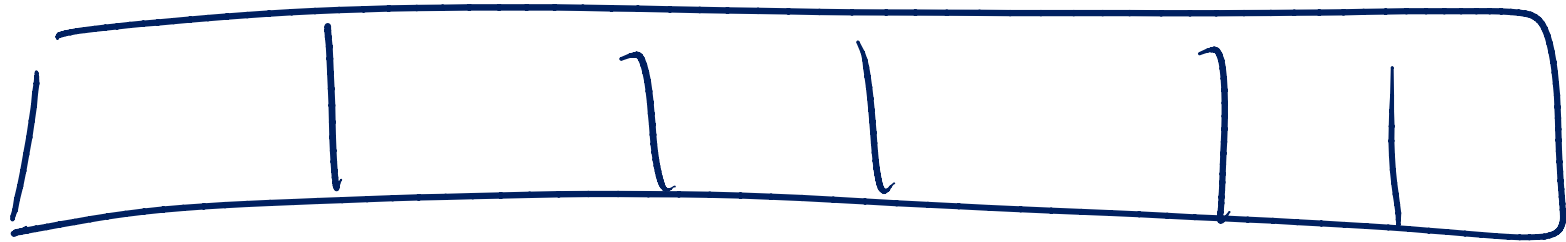
M1S14 When a record is deleted, it's marked as a ghost record. The ghost cleanup task does not physically delete the records – it removes the slot array entry that points to the record on the page. I.e. the record is unmarked as being a record and the area on the page is now free space – that just happens to have bits and bytes that used to be a valid record.

DBCC TRACEON

(662, -1)

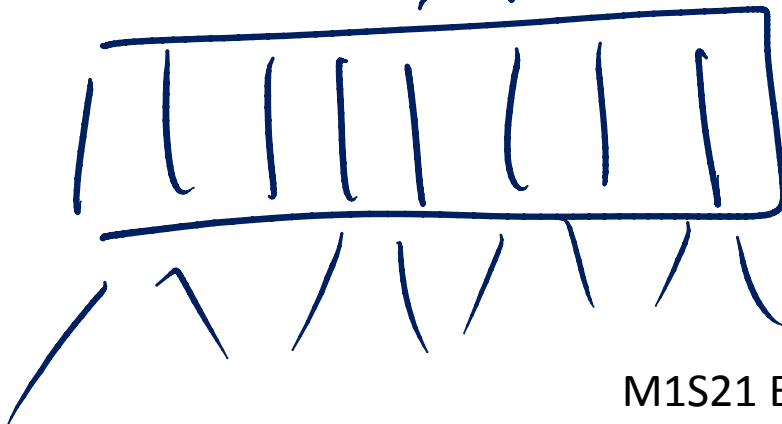


M1S14 For any trace flag that you want to effect background tasks, you must use the -1 in DBCC TRACEON to enable the trace flag globally instead of just for your connection.



0 - 7 8 - 15 16 - 23

MIXED



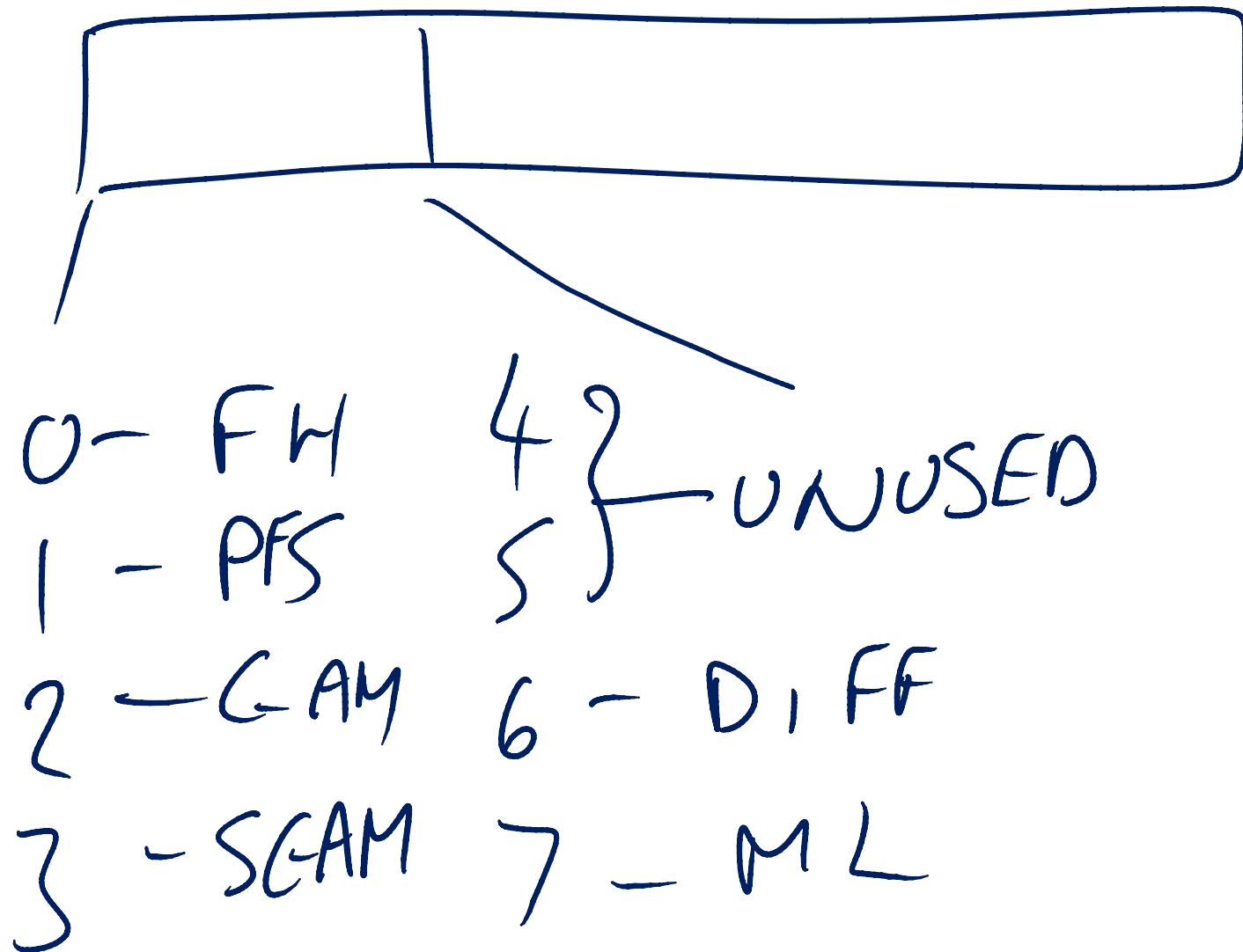
DEDICATED



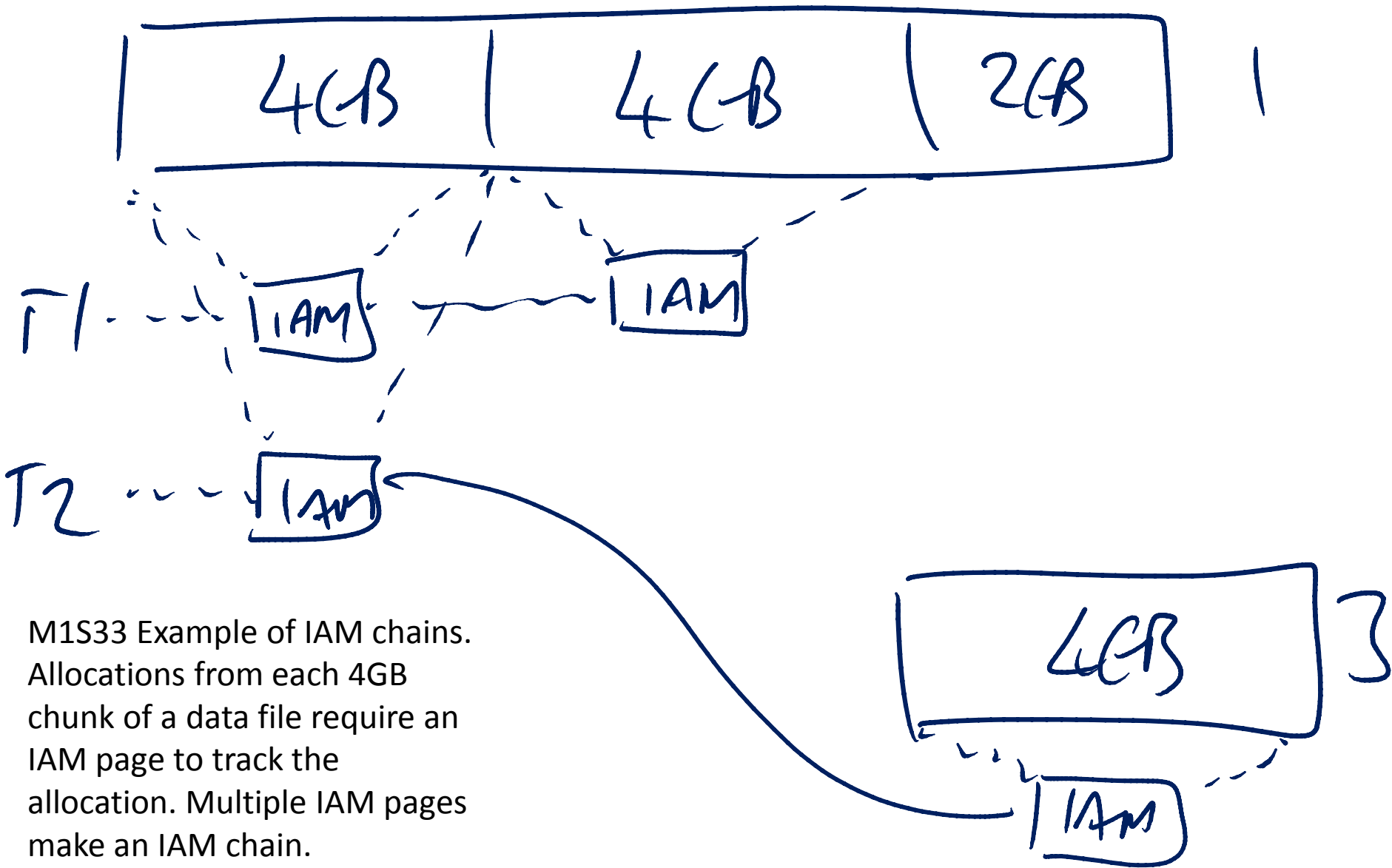
M1S21 Explaining the difference between mixed and dedicated extents. Mixed extent is one in which the 8 pages could be allocated to 8 different objects.

Dedicated extent is one where all 8 pages are reserved for exclusive use of a single object.

Later in the deck I refine 'object' to really be 'allocation unit'.



M1 The page types in the first extent in a data file

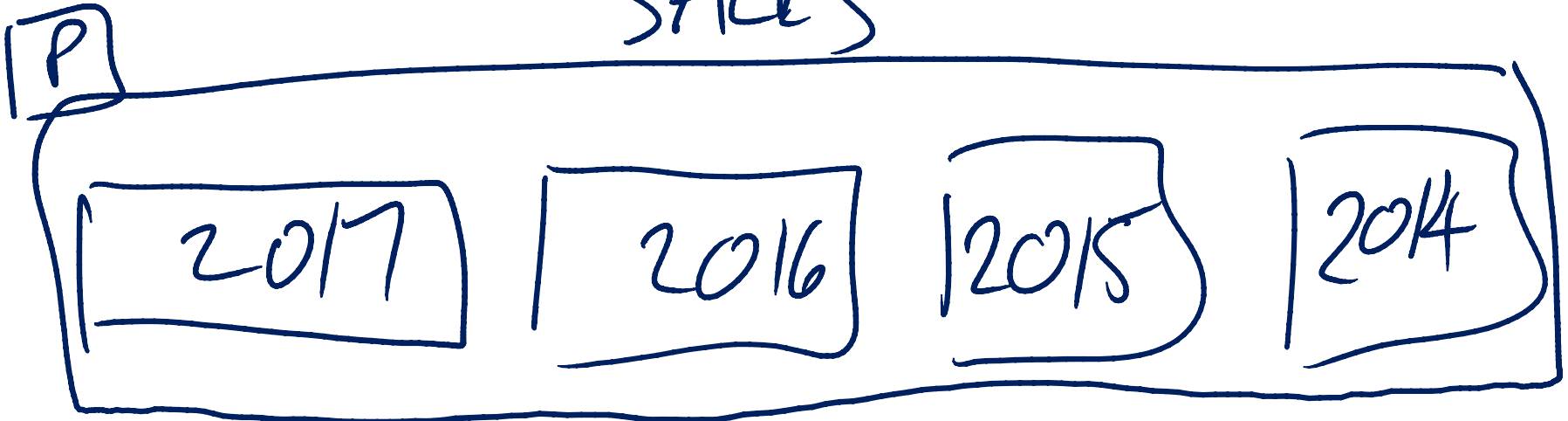


M1S33 Example of IAM chains. Allocations from each 4GB chunk of a data file require an IAM page to track the allocation. Multiple IAM pages make an IAM chain.

1TB SALES

M2S6 Splitting a large database up into multiple filegroups can enable you to do small, targeted restores or even a partial restore from scratch – in the event that the database is damaged or destroyed.

SALES



WA

OR

PARTS

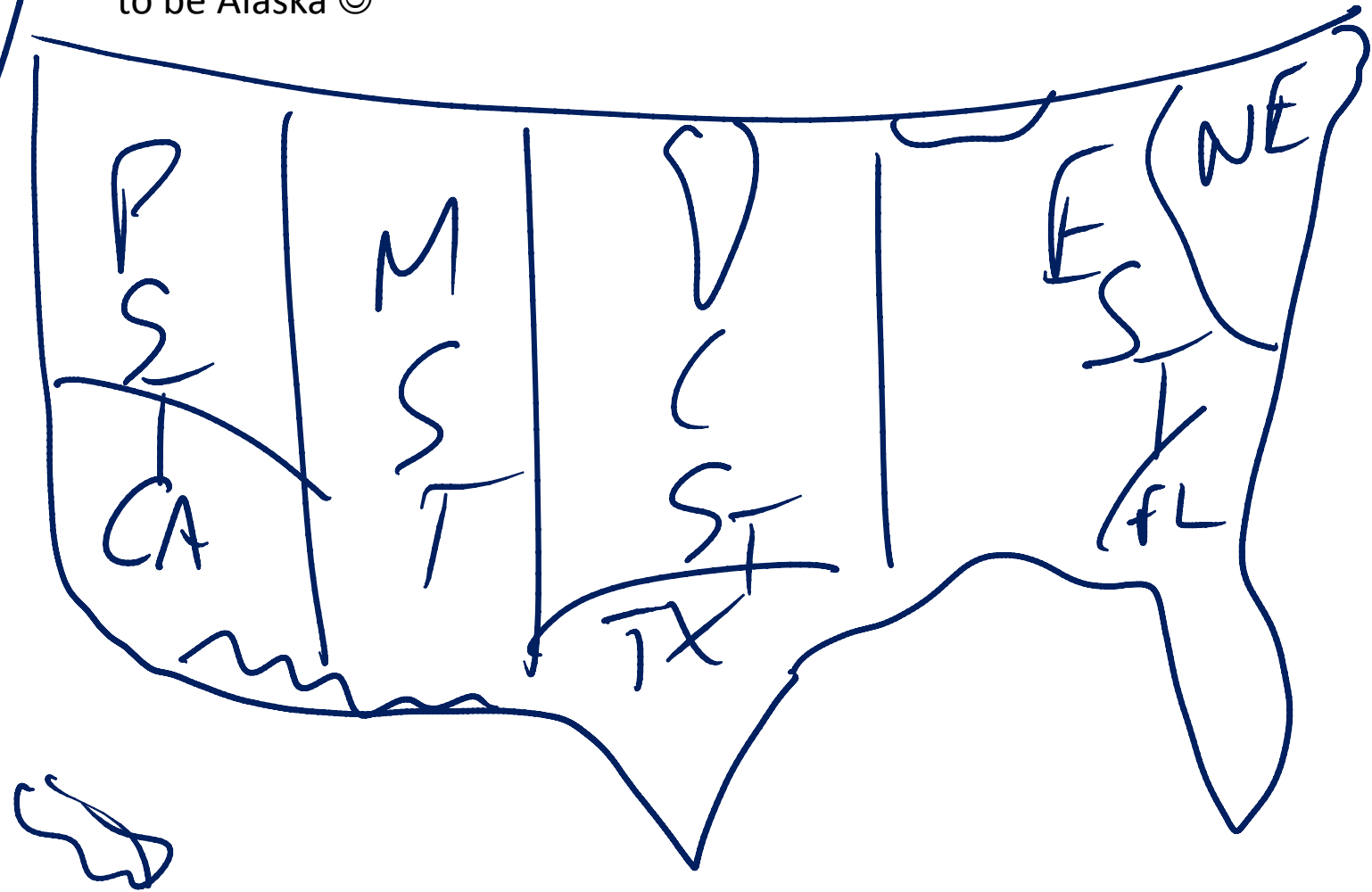
CA

ID

M2S6 Explaining some ways to split up a database to allow targeted restore. Consider time zone, population density, or other demographics.

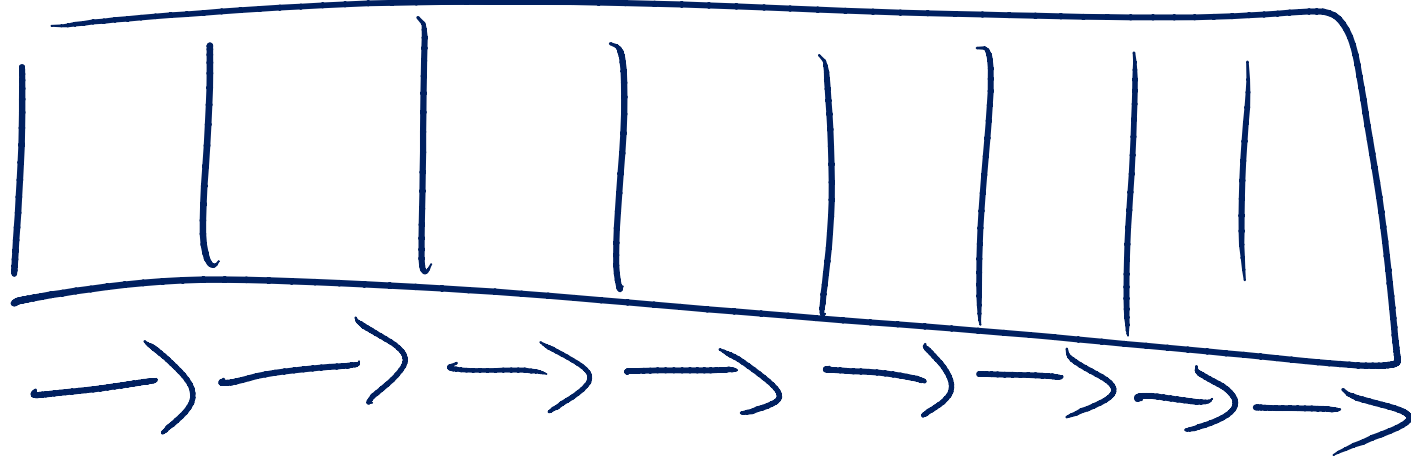
← This is
meant
to be Alaska ☺

7AM - 6PM



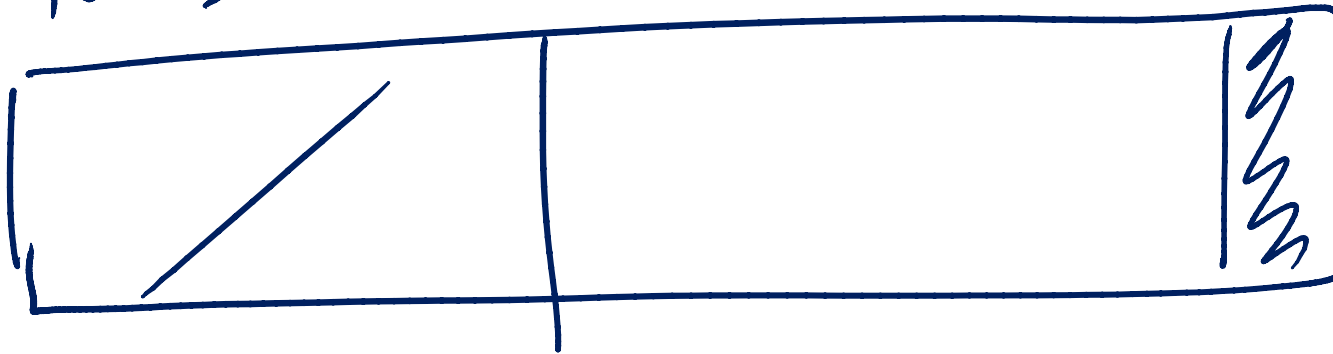
M2S6 Explaining some ways to split up a database to allow targeted restore. Consider time zone, population density, or other demographics.

L:

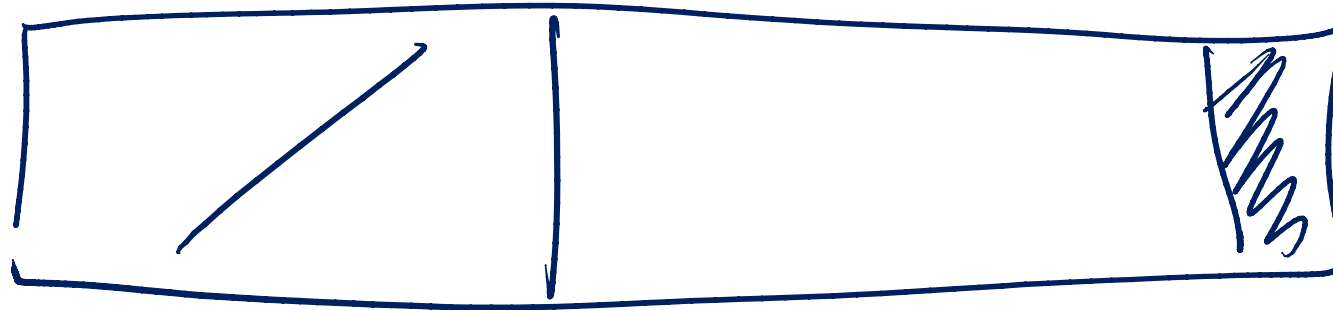


M2S10 If you put all your log files on one volume, its I/O characteristics become random. This may or may not become a performance problem, depending on the capabilities of the I/O subsystem.

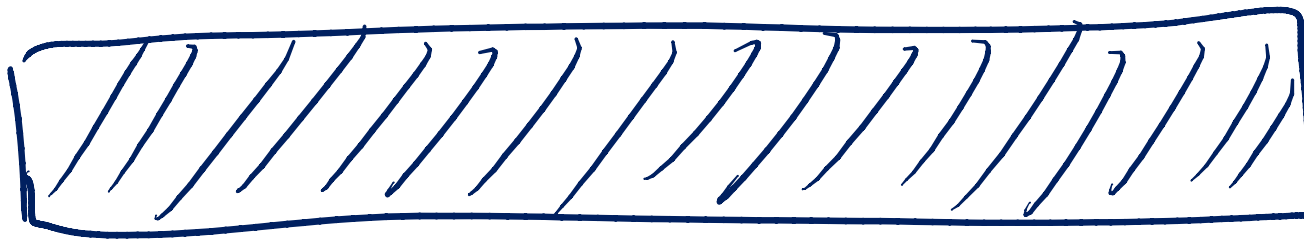
10GB



816085911



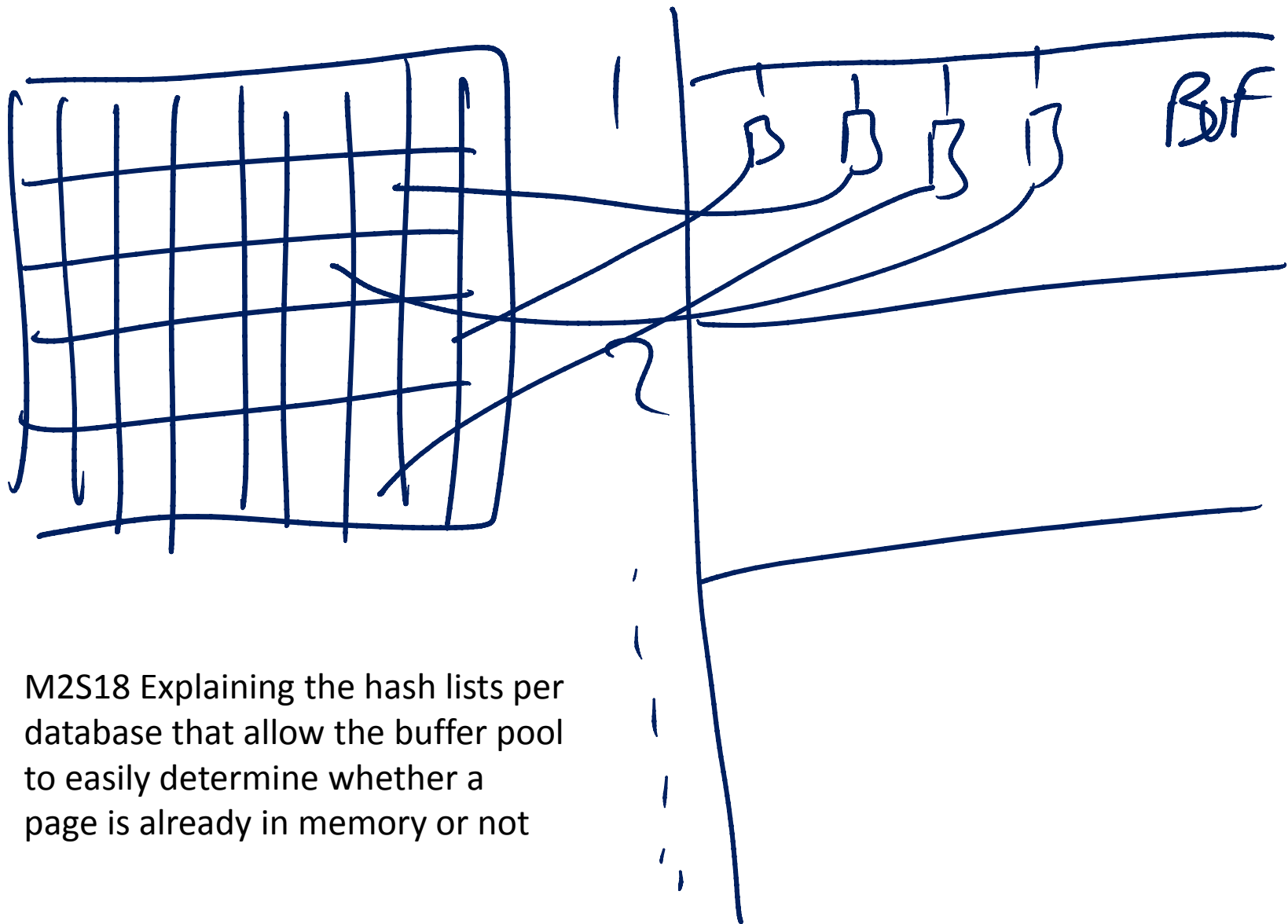
816085911



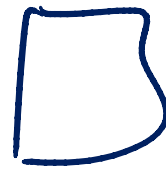
15870

M2S17 Explaining round-robin allocation and proportional fill. Each file has a proportional fill weighting. Allocations happen proportionally more frequently from files with more free space.

PROP FILL
WEIGHTING



M2S18 Explaining the hash lists per database that allow the buffer pool to easily determine whether a page is already in memory or not



LRU-K, K=2



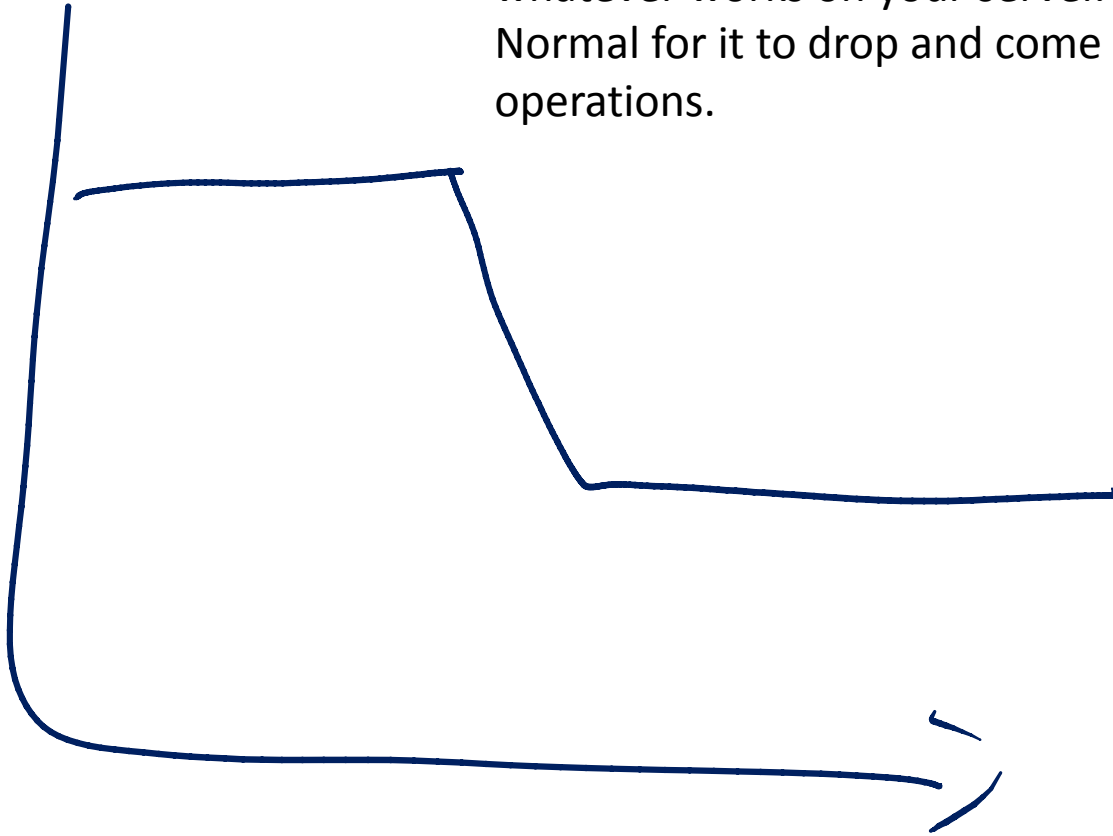
P.L.E.

LAZYWRITER

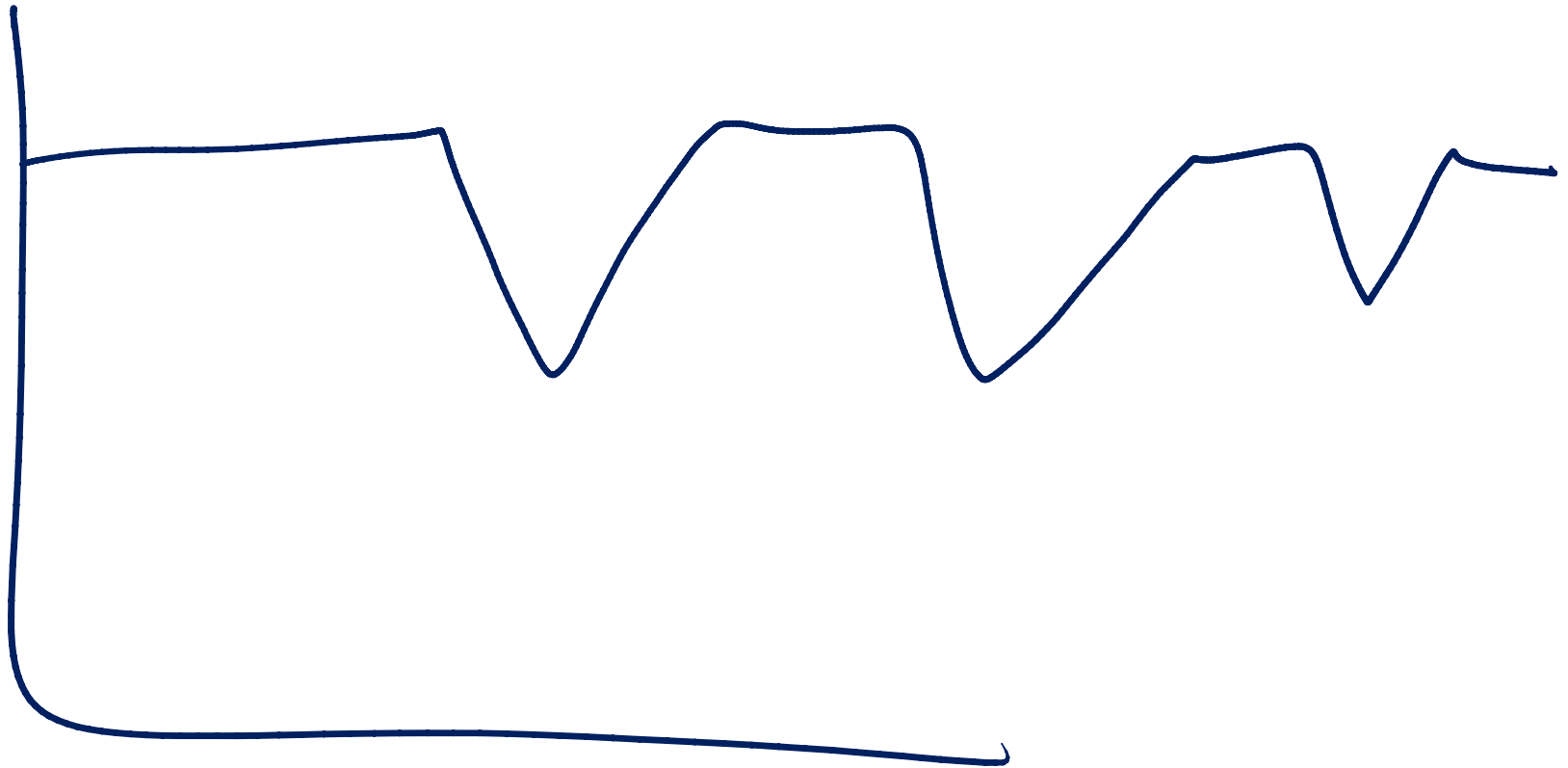
M2S18: The buffer pool implements a Least Recently Used algorithm, based on last two times the page was accessed. Done by the Lazy Writer background task.

P.L.E. ≤ 300

M2S18: PLE of 300 is an old, old threshold. PLE should be whatever works on your server. Problem if it drops and stays low. Normal for it to drop and come back up again after some operations.

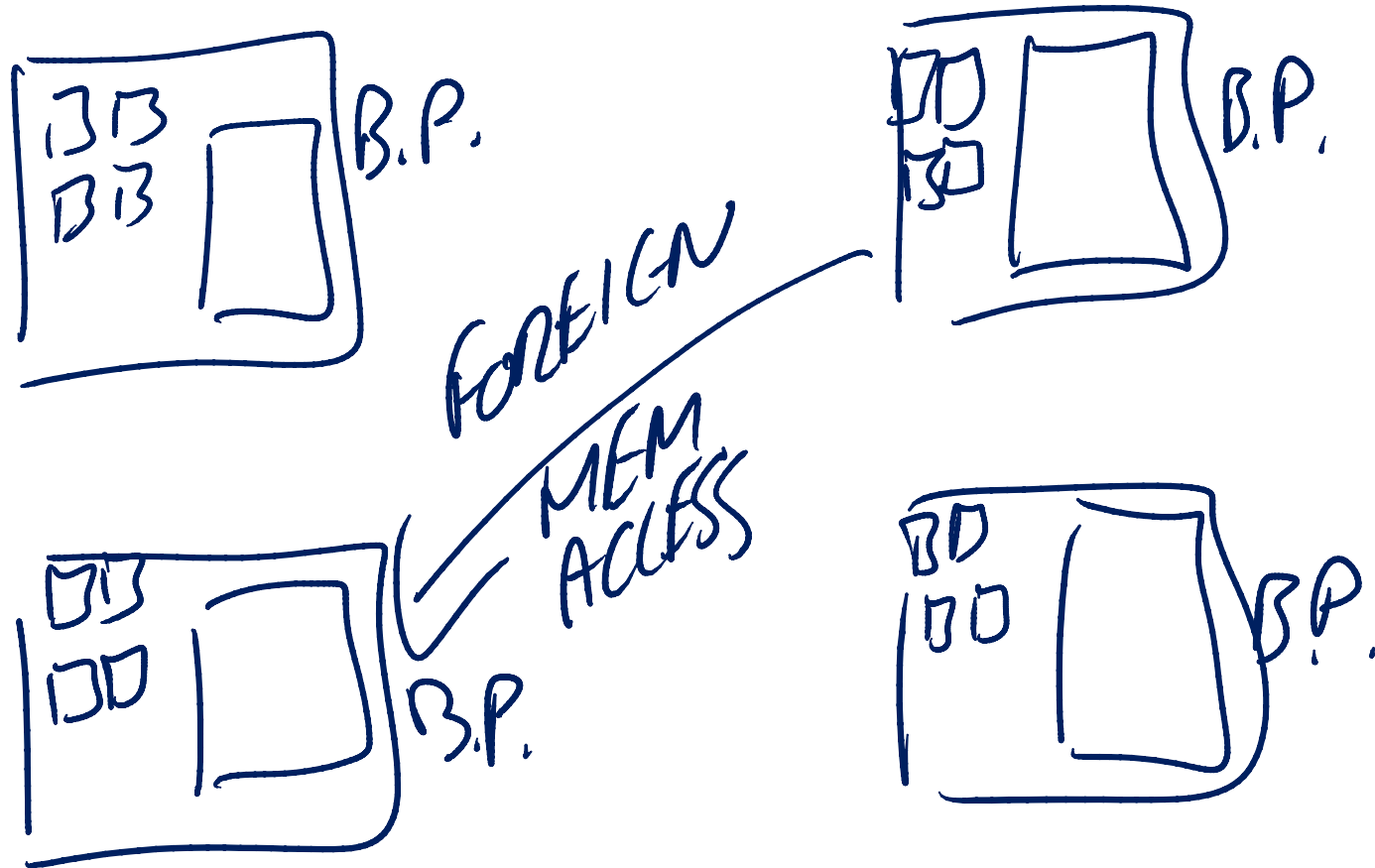


M2S18: PLE of 300 is an old, old threshold. PLE should be whatever works on your server. Problem if it drops and stays low. Normal for it to drop and come back up again after some operations.



NUMA

BUFFER NODE: PLE



M2S18: NUMA – Non Uniform Memory Access. Each processor (or group) has local memory, containing a partition of the buffer pool. 4 NUMA nodes = 4 equal-size partitions of the buffer pool.

NUMA - Sys. km.os-nodes

BUFFER MGR PLE.

X = HARMONIC MEAN

BUFFER NODE PLE

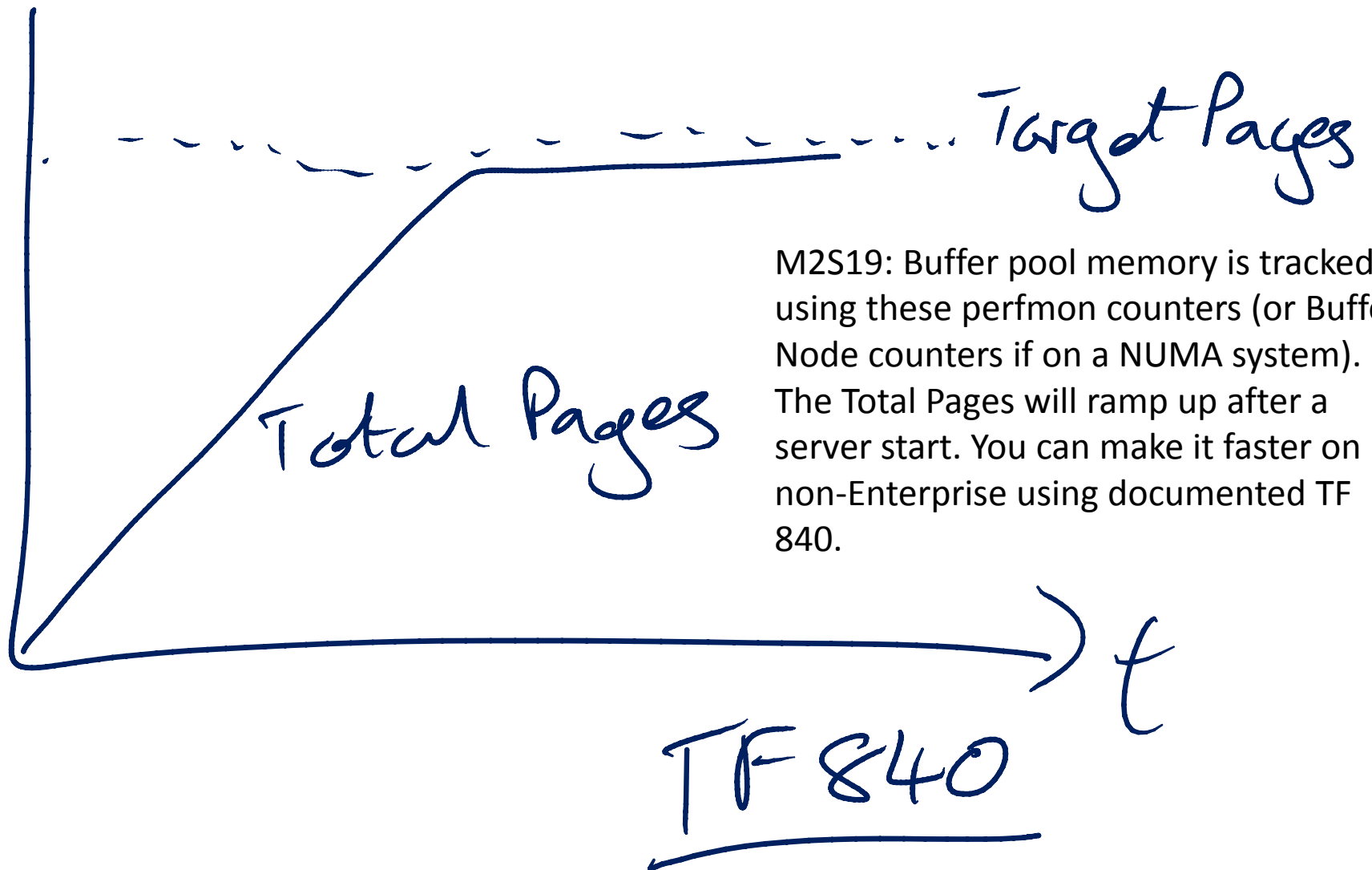
✓ M2S18: If you have NUMA, you need to monitor the PLE in each Buffer Node perfmon object, not the Buffer Manager PLE.

||
||
||

NUMA LINK

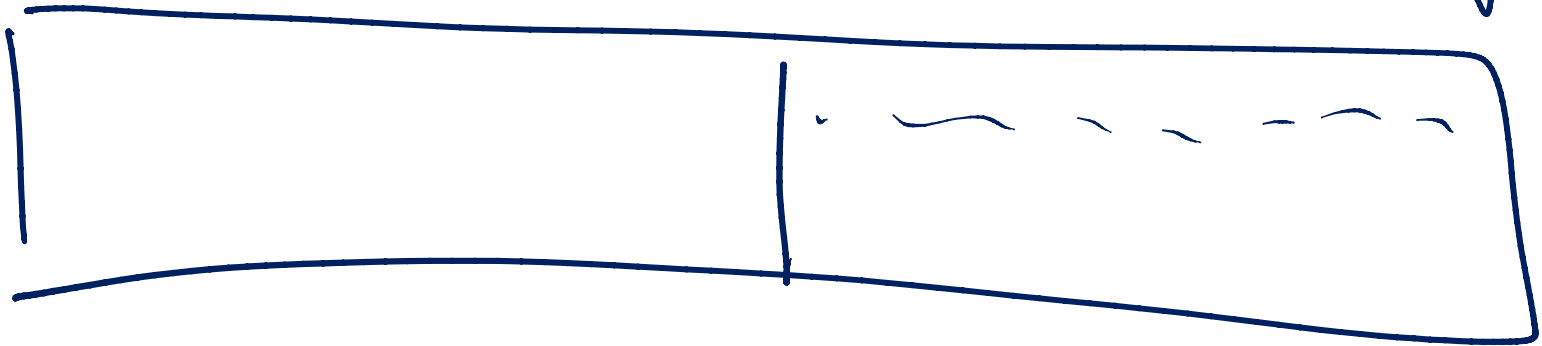
Someone asked for a link to info about configuring NUMA – best place to start is Jonathan's blog post category:

<https://www.sqlskills.com/blogs/jonathan/category/numa/>



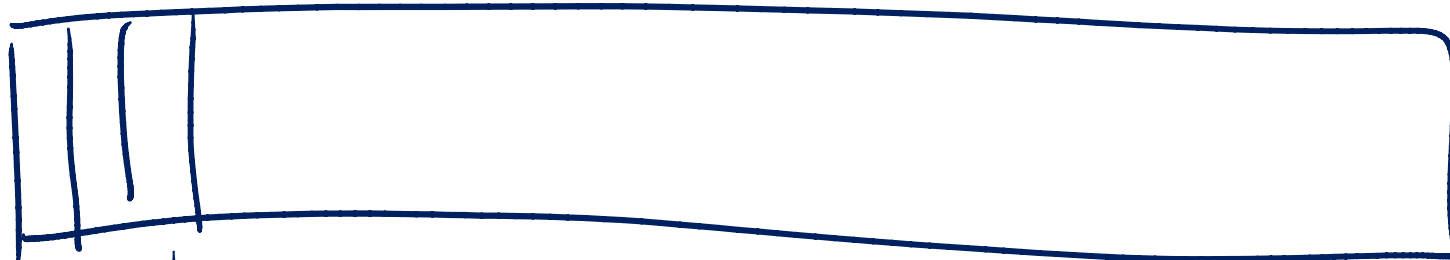
M2S19: Buffer pool memory is tracked using these perfmon counters (or Buffer Node counters if on a NUMA system). The Total Pages will ramp up after a server start. You can make it faster on non-Enterprise using documented TF 840.

SETFILEVALIDDATA



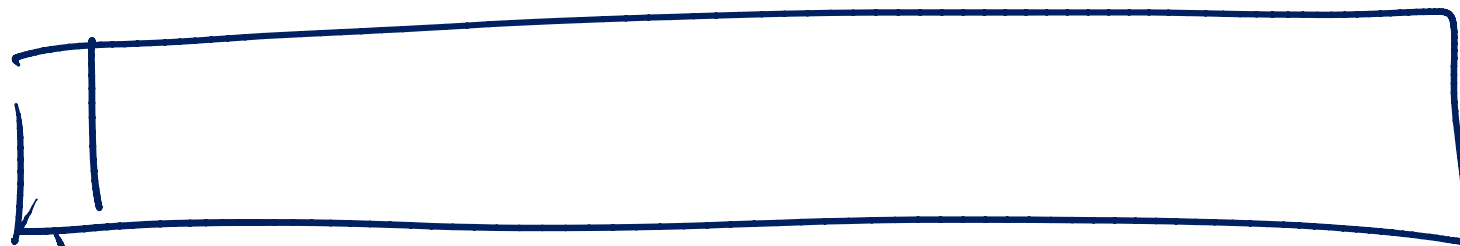
HIGH
WATER
MARK

M2S21 NTFS will not 'trust' a newly added portion of the file until the 'high water mark' of the file is advanced to include that section. This is usually done by writing zeroes to the file, but Instant File Initialization does it by calling the SetFileValidData API to do it 'instantaneously'.

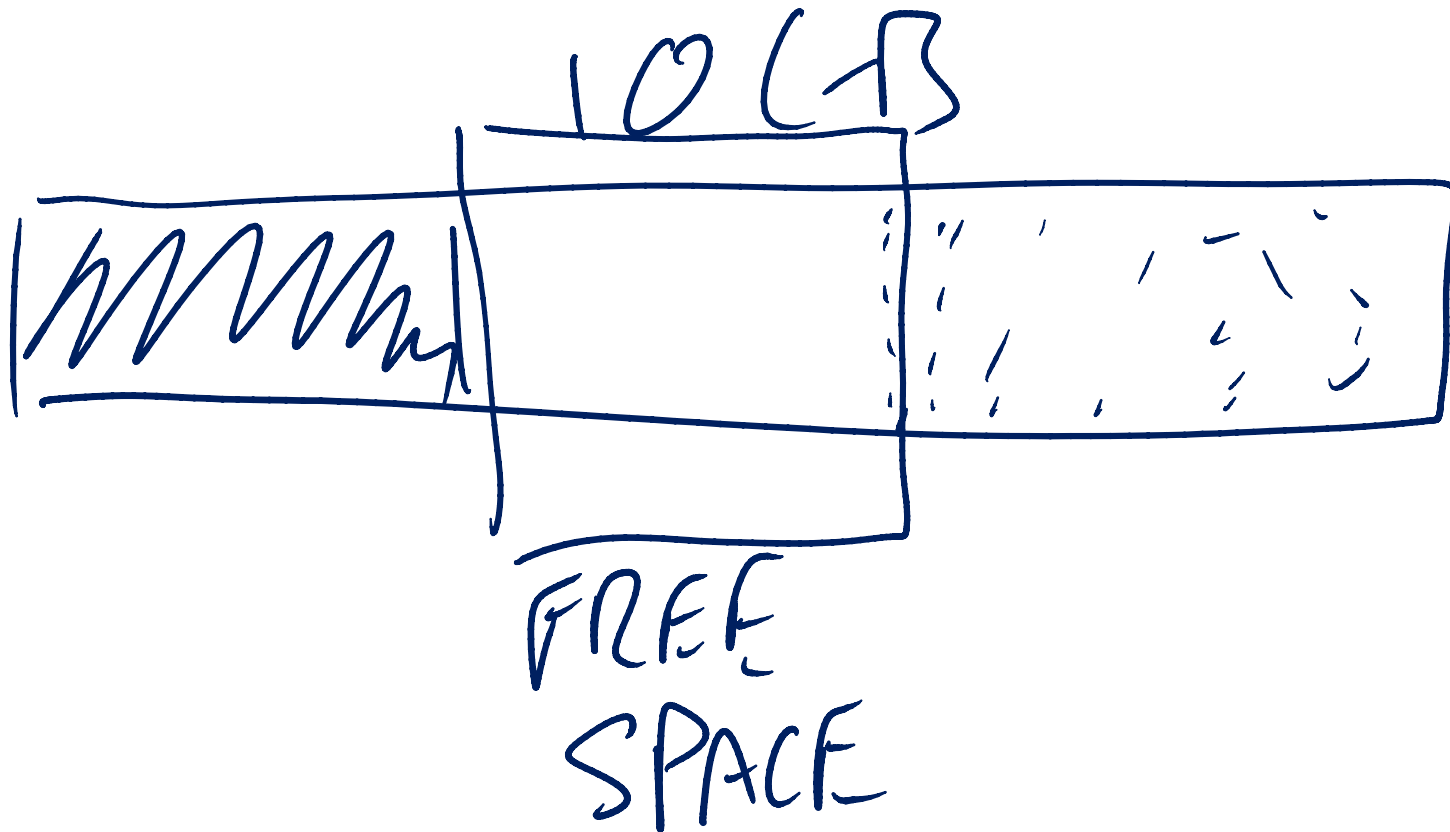


PFS
~~SGAM~~ T1118
2:1:1 ~~2:1:3~~

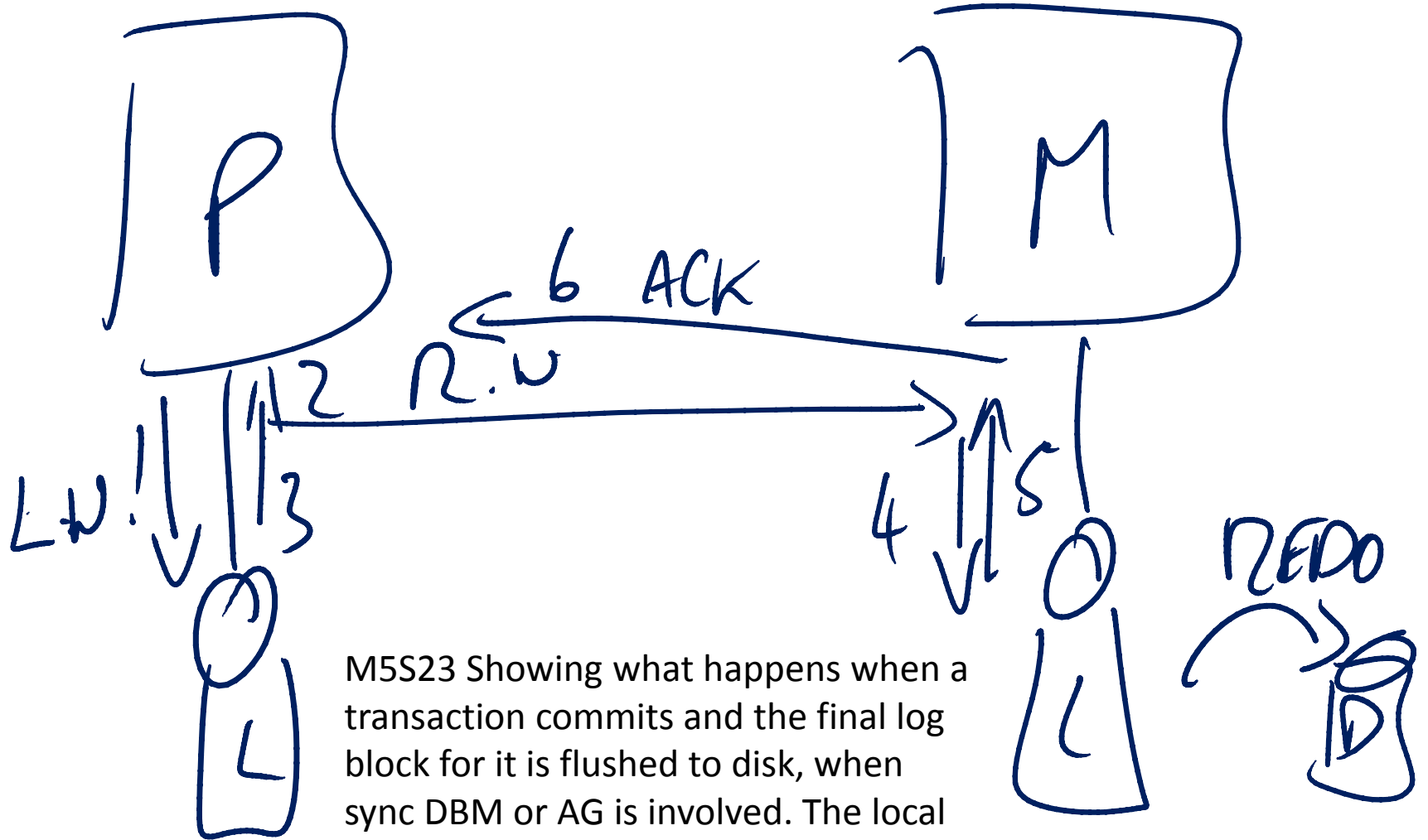
M2S51 Explaining PFS and SGAM contention in tempdb and how adding multiple files spreads the contention over multiple files, thus reducing it.



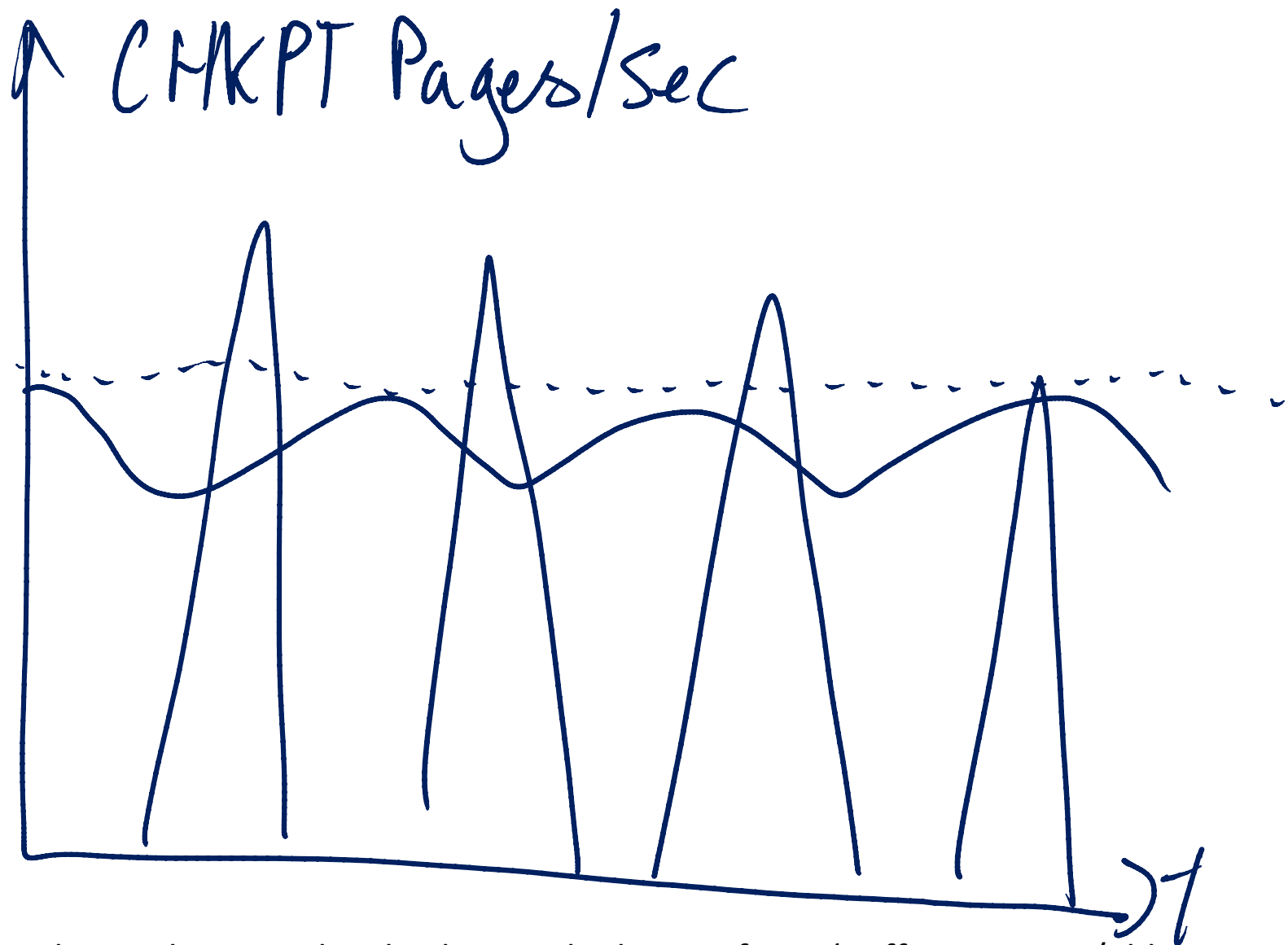
PFS 2:3:1



M5S14 Discussing log space reservation. On the left are the log records, on the right is the reserved space (conceptually). The bit in the middle isn't needed, so the file could be shrunk that much, but it's not possible to shrink the file such that there's not enough reserved space.

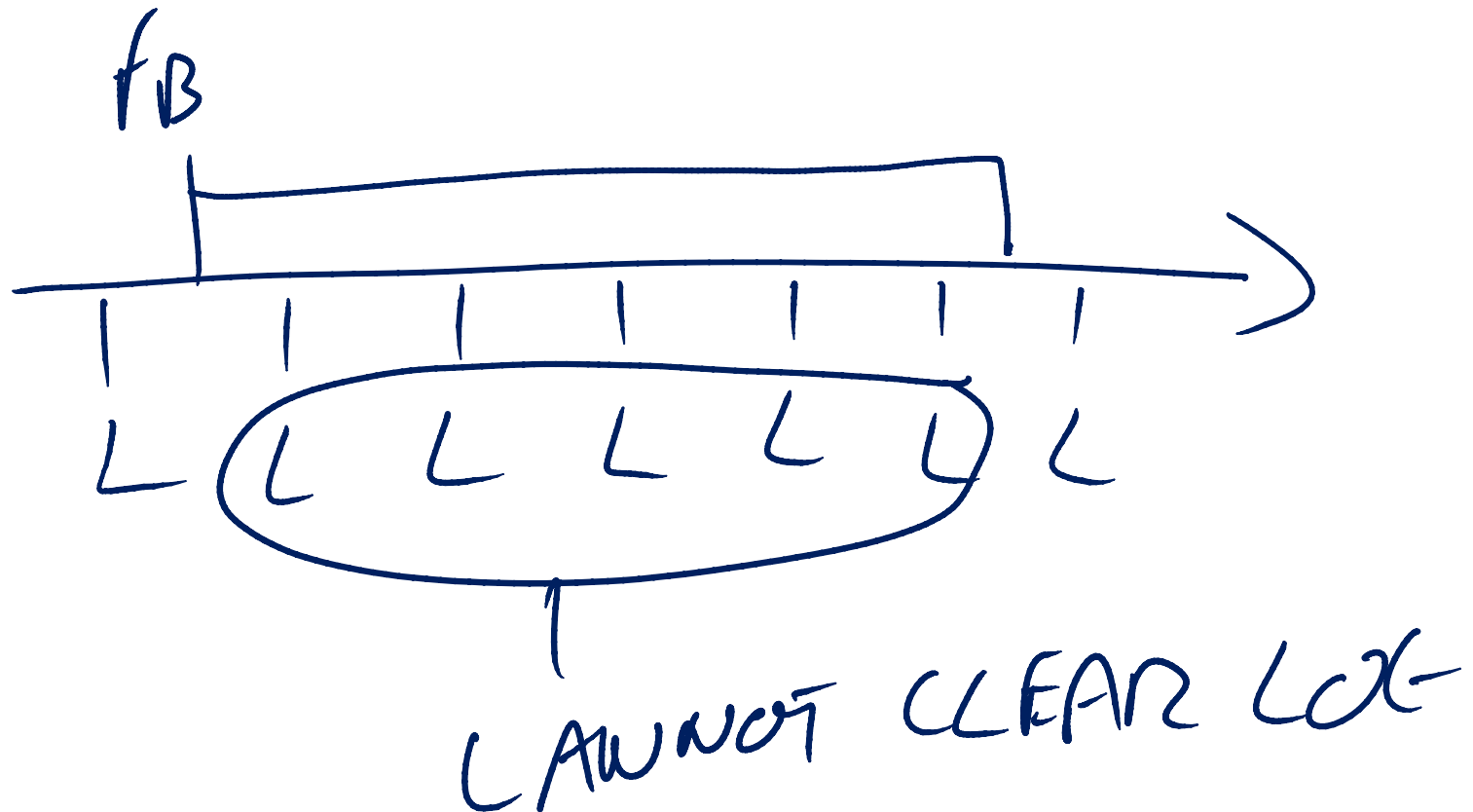


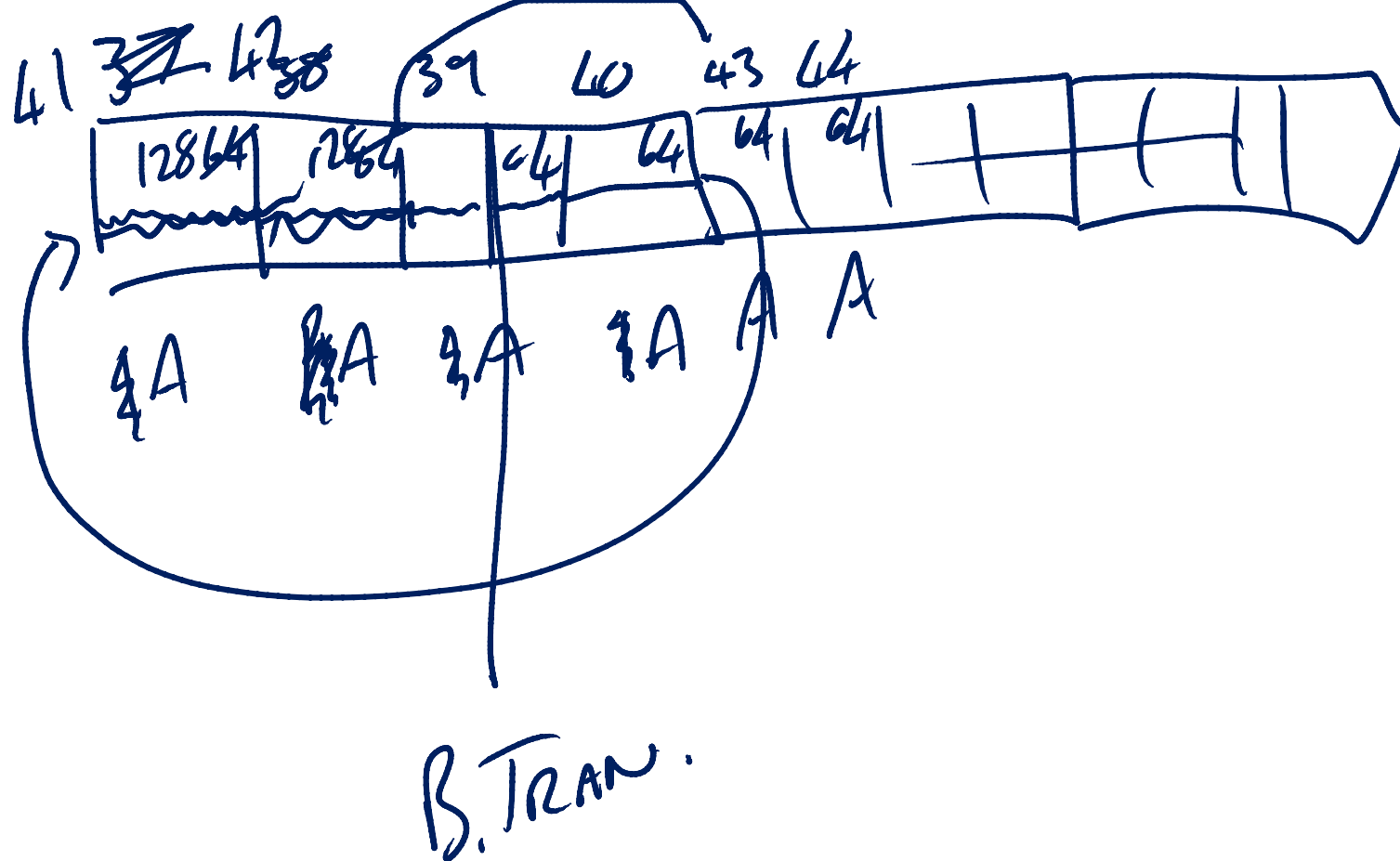
M5S23 Showing what happens when a transaction commits and the final log block for it is flushed to disk, when sync DBM or AG is involved. The local I/O will complete first (3) but the transaction commit cannot complete until the mirror acknowledges the remote I/O (6).



M5S35 Explaining how regular checkpoints look in perfmon (Buffer Manager/Chkpt pages/sec counter) and how they may exceed I/O capacity (dotted line). More frequent manual checkpoints, or using indirect checkpoints in SQL 2012+, can mean more constant, lower IO load.

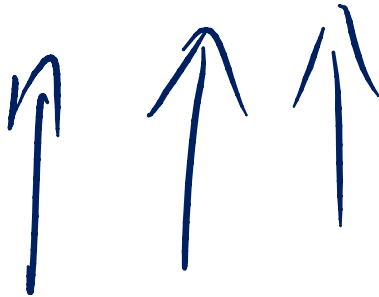
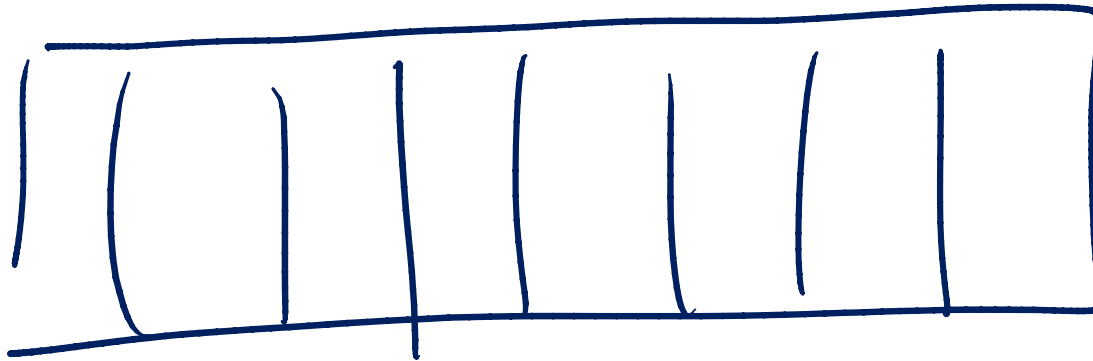
M5S47 Explaining how concurrent log backups while a data backup is running cannot clear the log. When the data backup completes, deferred log clearing takes place. This gives the impression that a data backup clears the log.





M5S53 Explaining the circular log demo. We discussed log space reservation that happens to ensure the active transactions can roll back without having to grow the log – this accounts for extra log growths. Space reserved in the log does not make a VLF become active – as there are no log records written out, just empty space reserved.

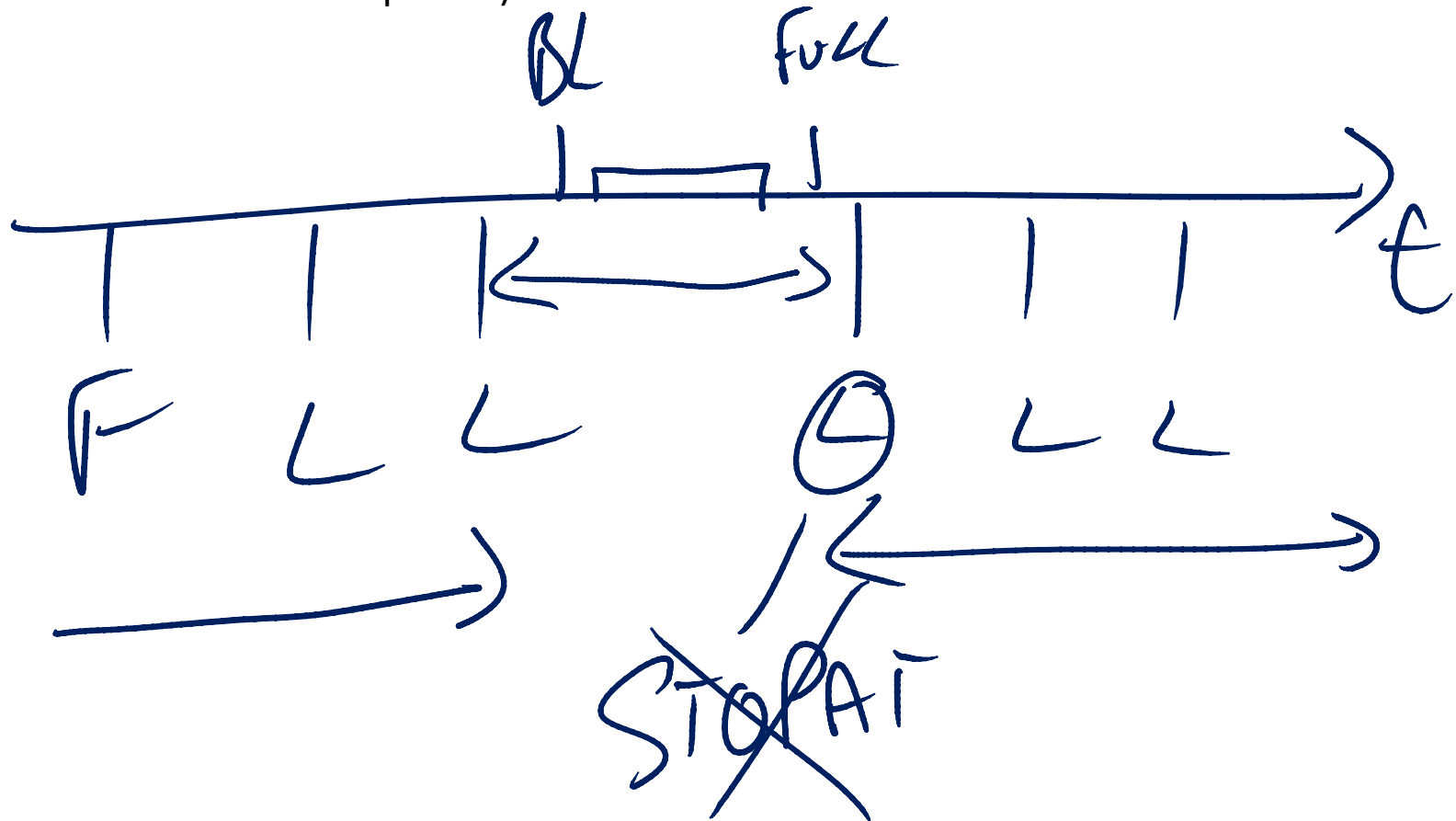
EXTENT X



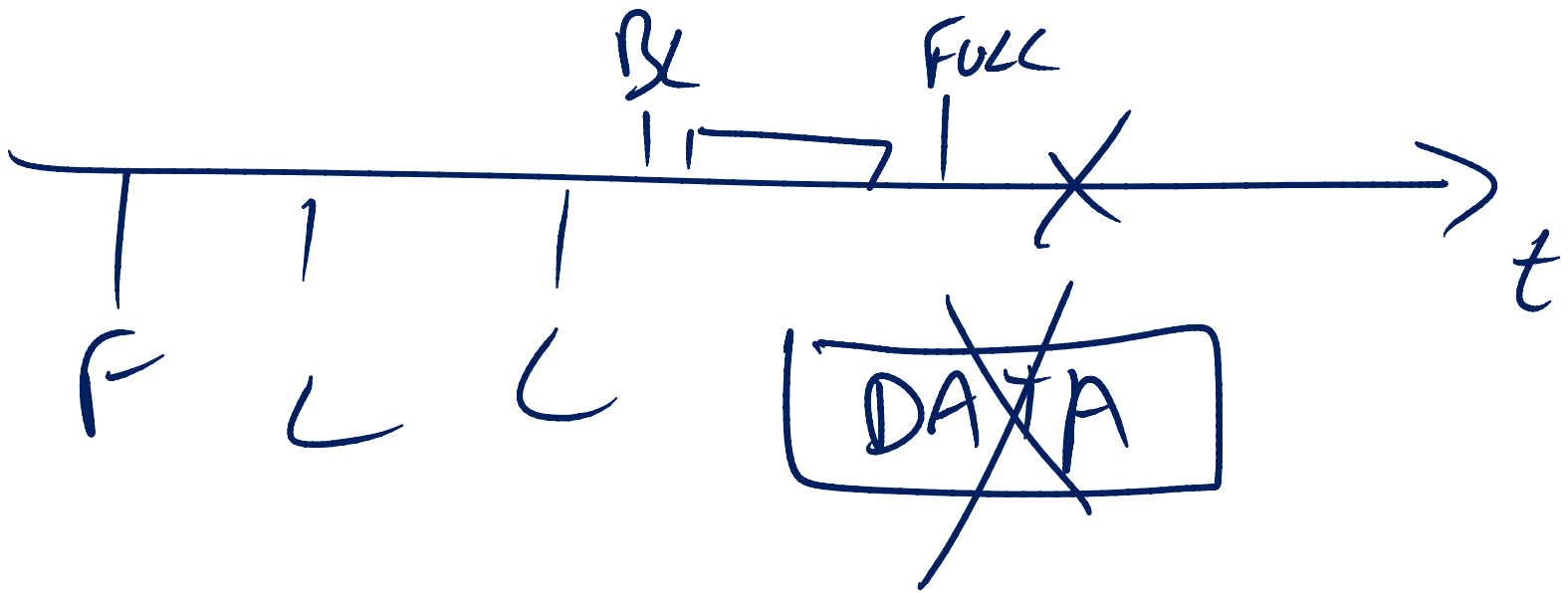
PROBE

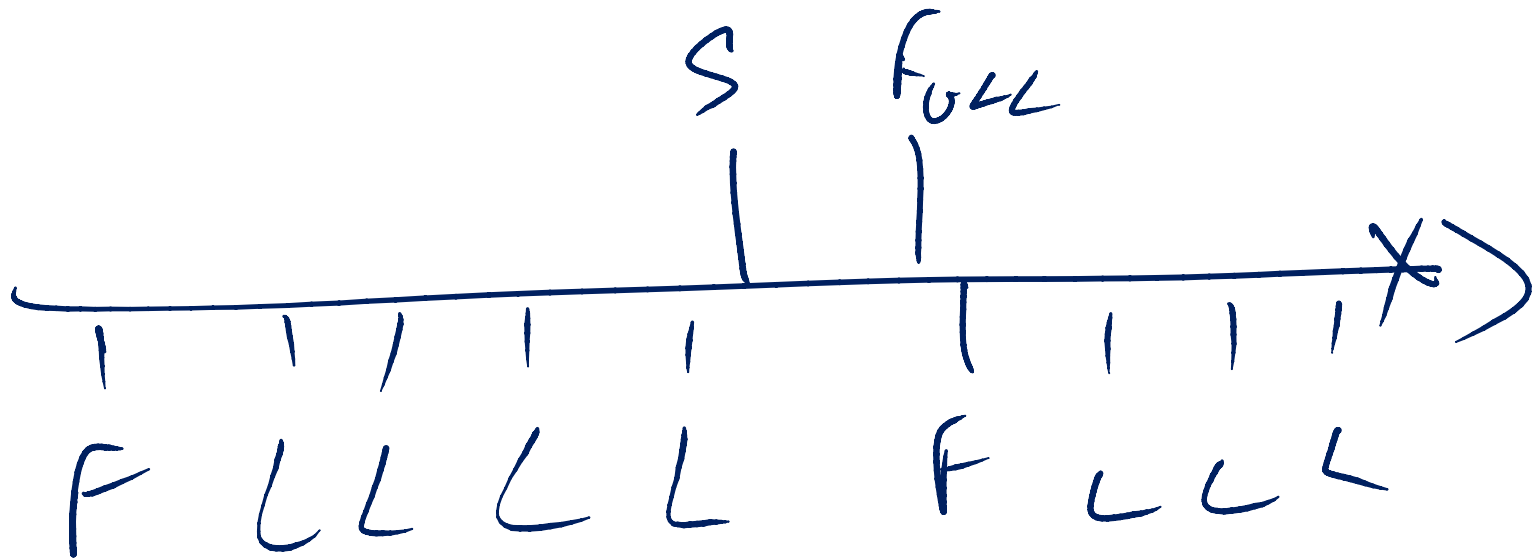
M5S68 Explaining part of the deferred-drop functionality. Before an extent can be deallocated, an X lock is acquired on it and all 8 page locks are probed (instant acquire X lock and release) to make sure nothing else has them locked.

M5S69 Explaining that if a minimally-logged operation is present in a log backup, you cannot do a point-in-time restore to any time covered by that backup (except the start or end of that time period).

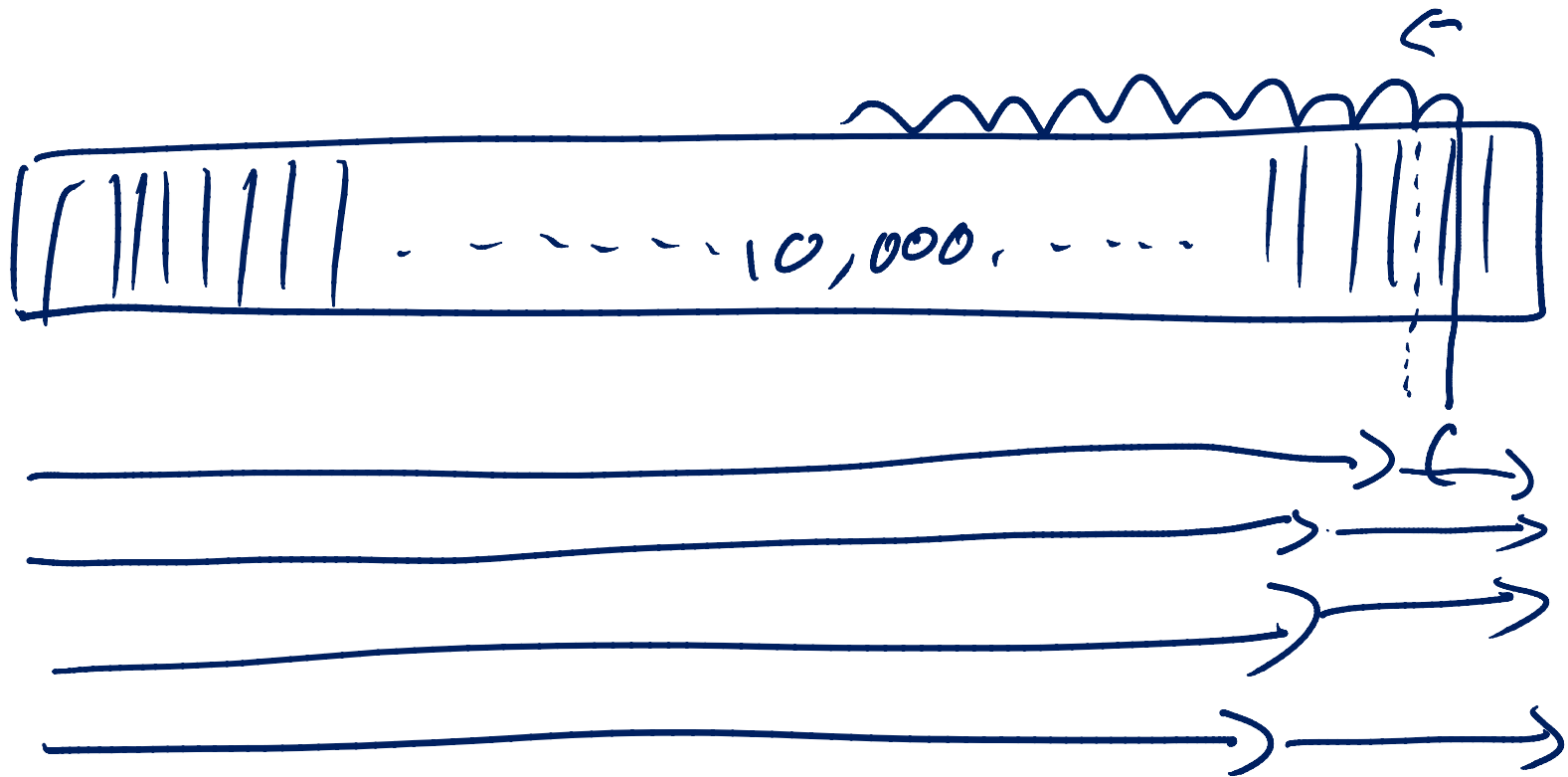


M5S69 You cannot do a tail-of-the-log backup after a crash that destroys the data files if there has been a minimally-logged operation since the last log backup. Well in 2008 R2+, you can, but the log backup is invalid.

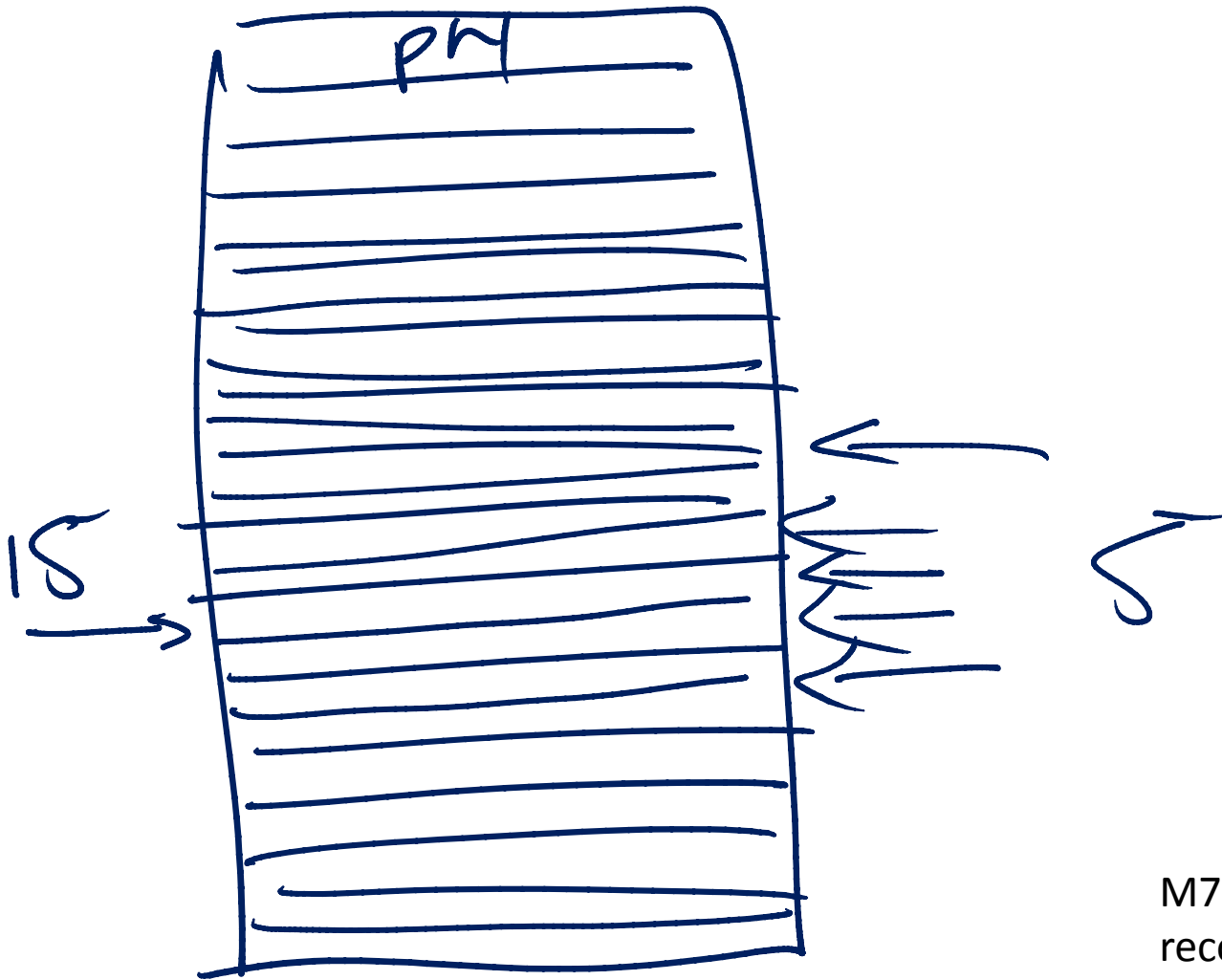




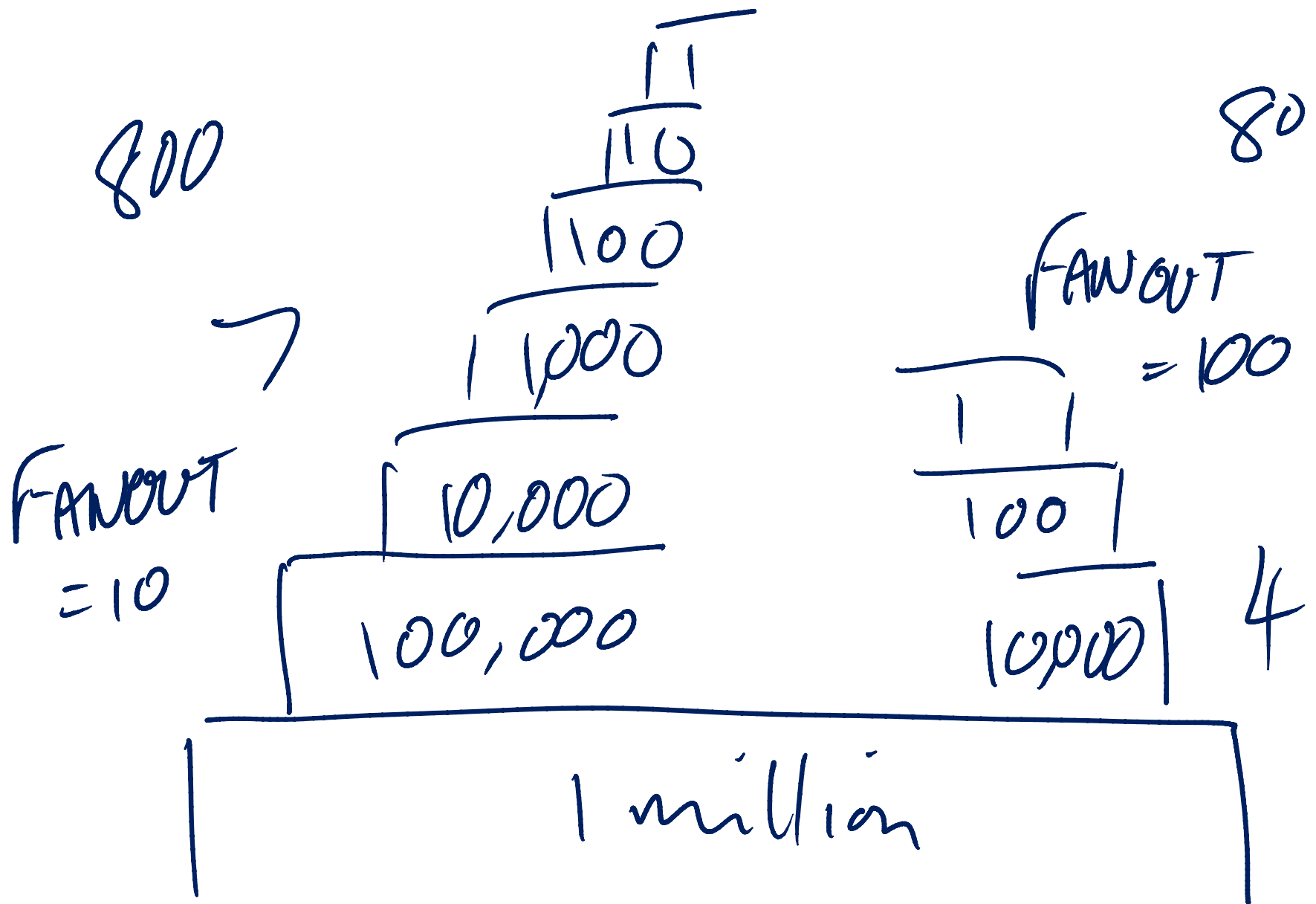
M5S70 If you ever have to switch to the simple recovery model, when you switch back to full, reestablish the log backup chain using a data backup. You can use a differential backup to do this.



M5S71 VLF fragmentation is bad because of having to search through the VLFs for a particular VLF sequence number. See the KB article link on the slide for more info.



M7S4 Explaining how a record is found on a page using binary search rather than brute force search.



M7S5 Explaining how a larger index key size reduces fanout, thus requiring more levels in the b-tree. Each level incurs the cost of a binary search, so this does matter.

COMMENTS MEMBERID

PAUL	ERIK	GLENN	TIM
------	------	-------	-----

KIMBERLY

M7S19 The MySpace story. The tiny portion is for the comments on Kimberly's page because she wasn't very popular back then... 😊

BEGIN TRAN
INSERT → SYSTEM TRAN

M7S22 A page split is done by a system transaction that commits and does not ever roll back.

SPLIT

FINISH ——— COMMITS
COMMIT TRAN

M7S28 The numbers from the page split logging cost demo.

	NO SPLIT	SPLIT	?
1000	1256	7580	$\times 5$
100	356	6772	$\times 18$
10	268	11172	$\times 45$



$$100 \times 40$$

$$50 \times 40$$

$$25 \times 40$$

$$12 \times 40$$

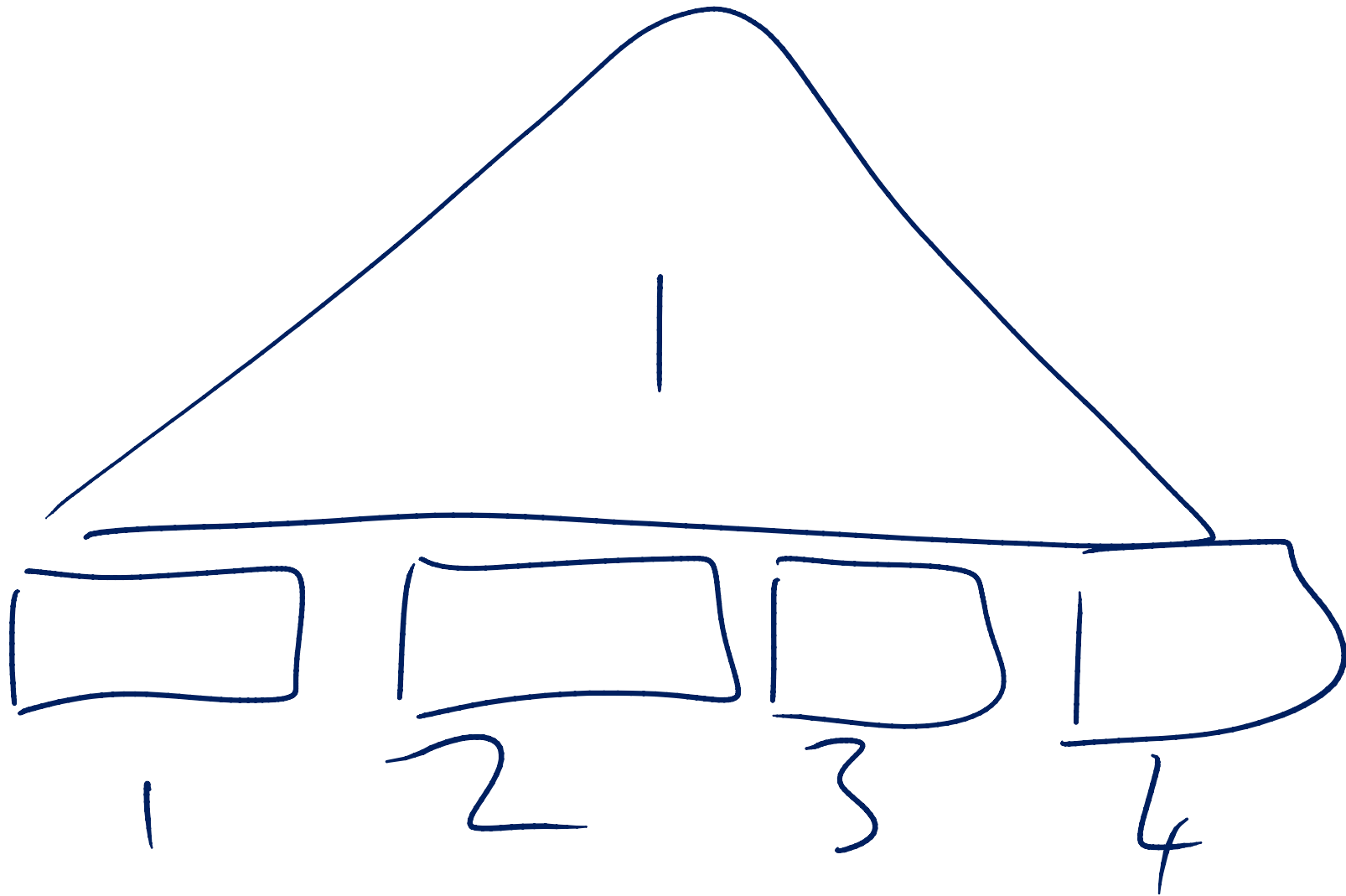
$$6 \times 40$$

$$3 \times 40$$

$$1 \times 40$$

8000

M7S22 Page splits aren't always smart. Sometimes it will split and there still isn't enough space, so it splits again, and again, and again....



M7S38 Discussing how DOP affects contiguity. With DOP 4, 4 threads are building portions of the new index in parallel. They will all be allocating extents, leading to extent fragmentation (and hence logical fragmentation).



M7S44 Explaining staggered index maintenance with database mirroring. See <http://bit.ly/IS04W5> for more explanation